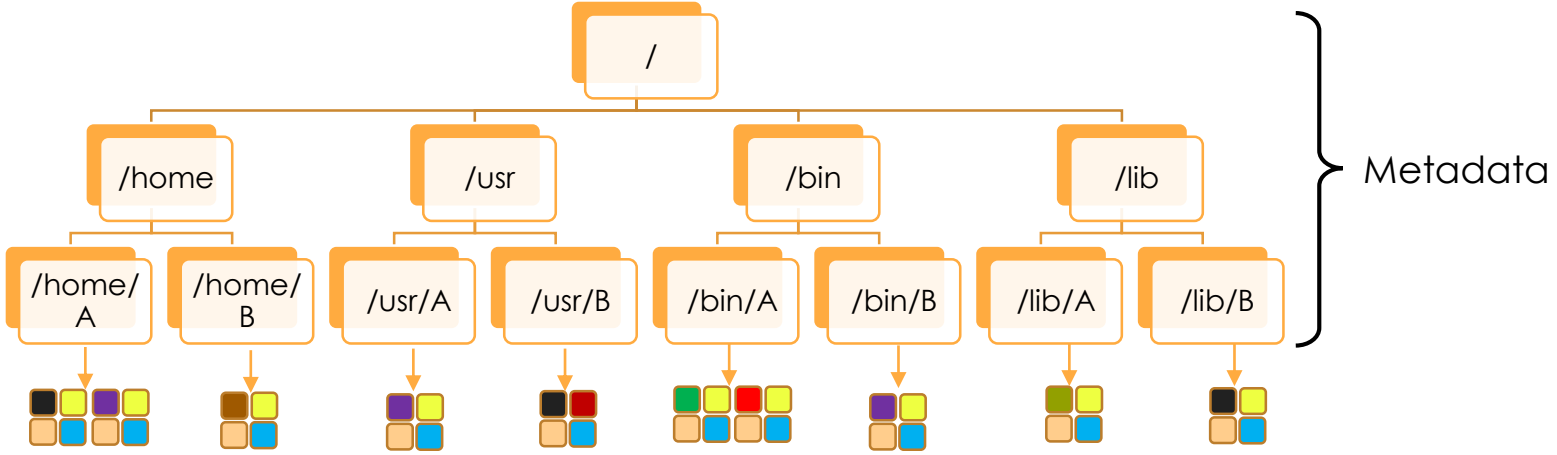


HopsFS

Scaling Distributed Hierarchical File Systems
Using NewSQL Databases

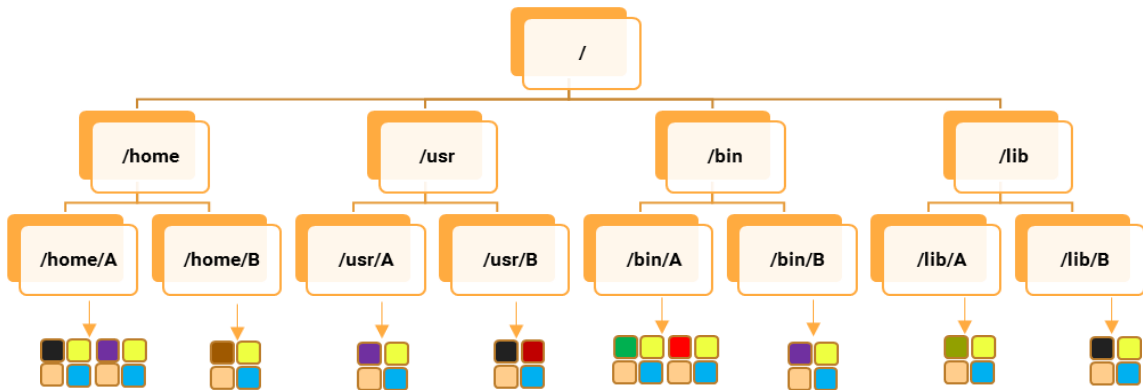
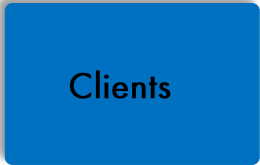
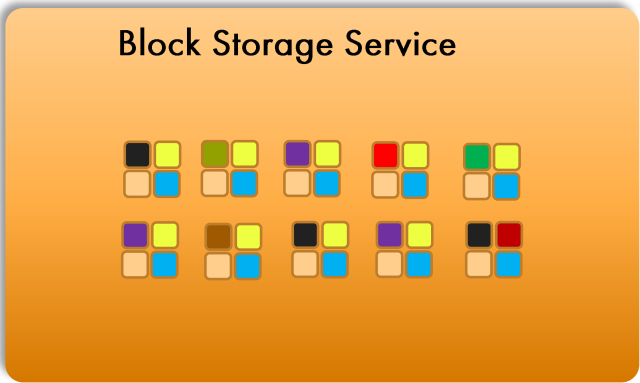
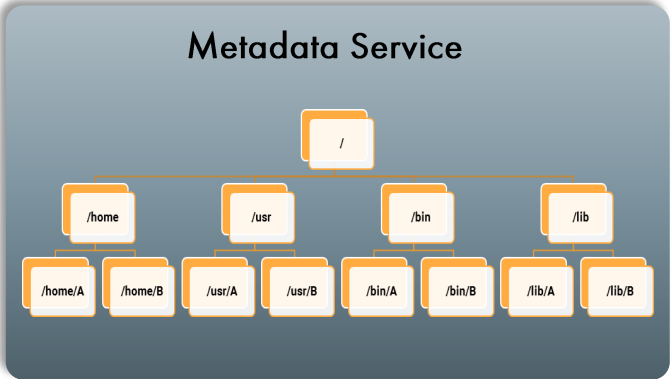
Salman Niazi

Hierarchical File System



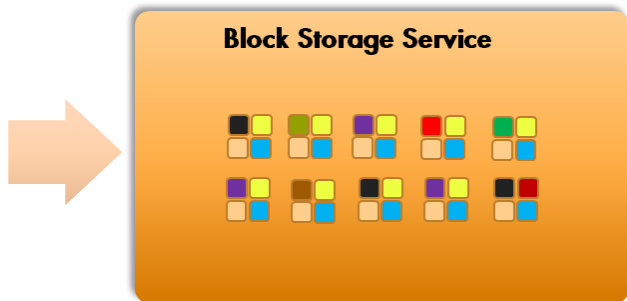
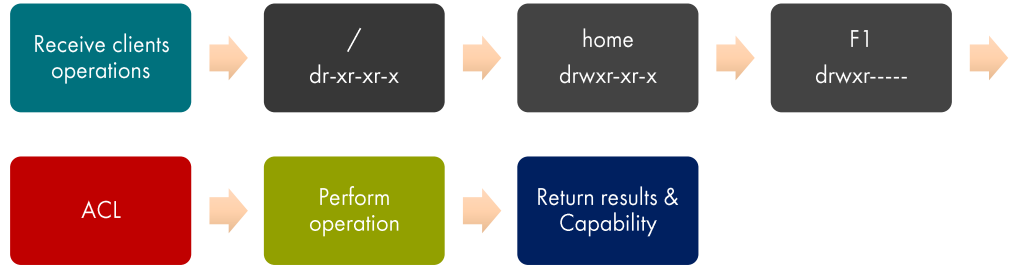
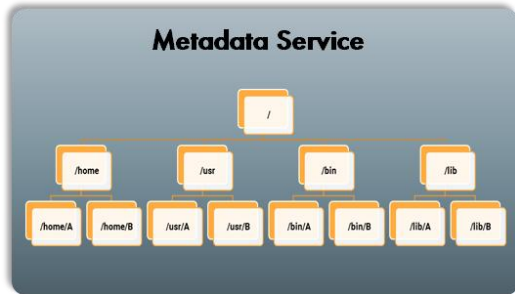
- Strongly consistent metadata.
 - Atomic file system operations, such as, move and create

Distributed Hierarchical File System



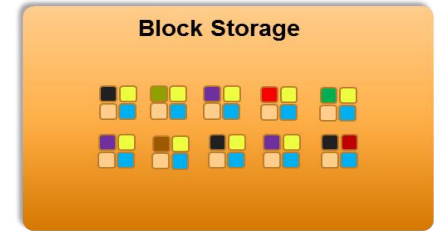
Typical Hierarchical File System Operation

- {operation} [flags] {path(s)}
- cat /home/F1



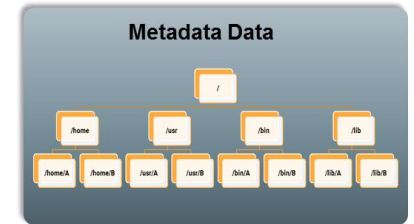
Data Blocks Storage Layer

- Thousands of servers
- Uses data replication and erasure coding for high availability



Metadata Service Layer

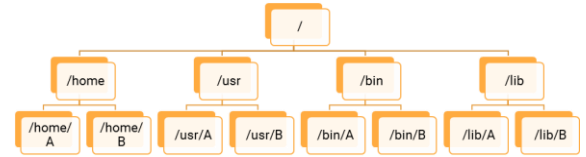
- Atomic File System Operations
- Due to complexity of metadata service monolithic architecture is the most popular solution
 - HDFS, GFS, AFS



Why not use Databases?

- Metadata consists of lots of very small data
- Databases specialize in storing and manipulating large amounts of small data

- Traditional databases do not provide high throughput required by distributed file systems
- High operational latencies for resolving file paths



Using Databases

- WinFS by Microsoft
- GiraffaFS, CassandraFS, CalvinFS



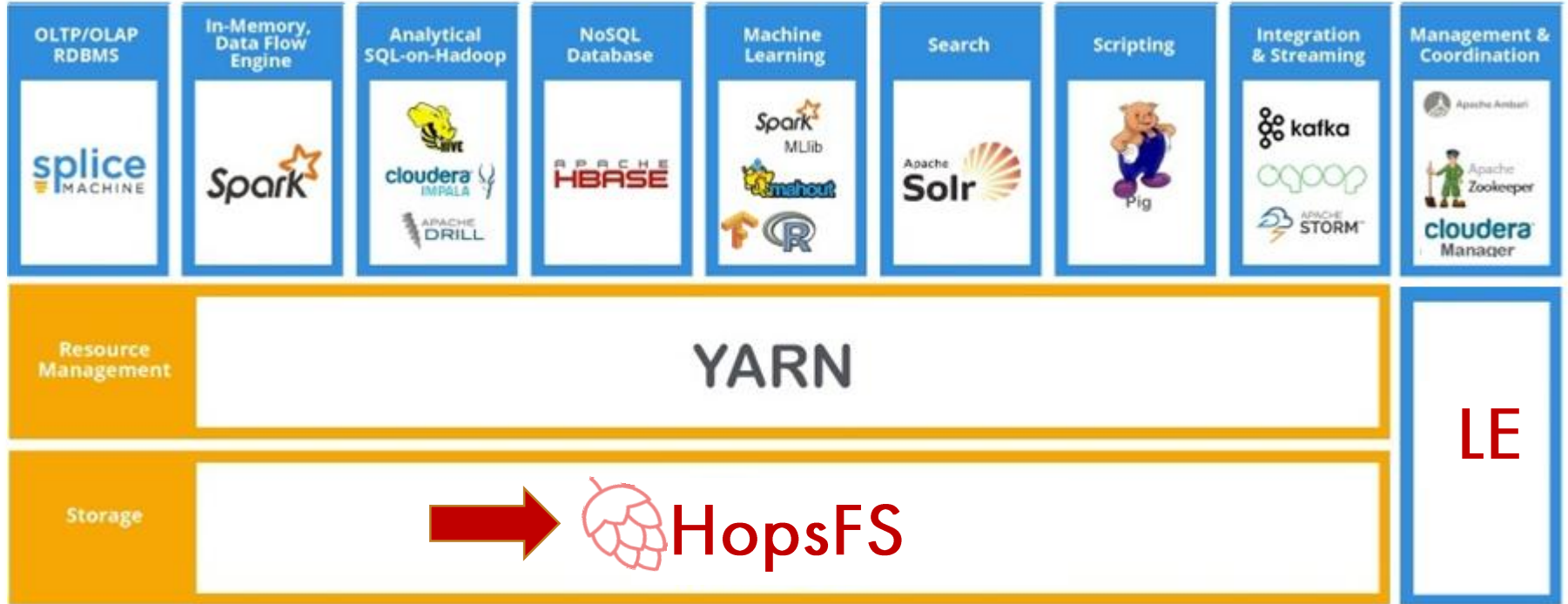
inode	...
/	
/usr	
/usr/F1	
/usr/F2	
/usr/F3	

Denormalized

ID	name	PID	
0	/	-1	
1	usr	0	
2	F1	1	
3	F2	1	
4	F3	1	

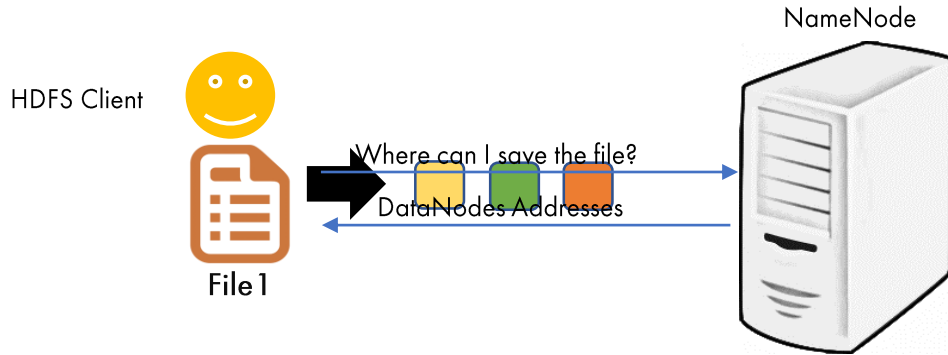
Normalized

HopsFS



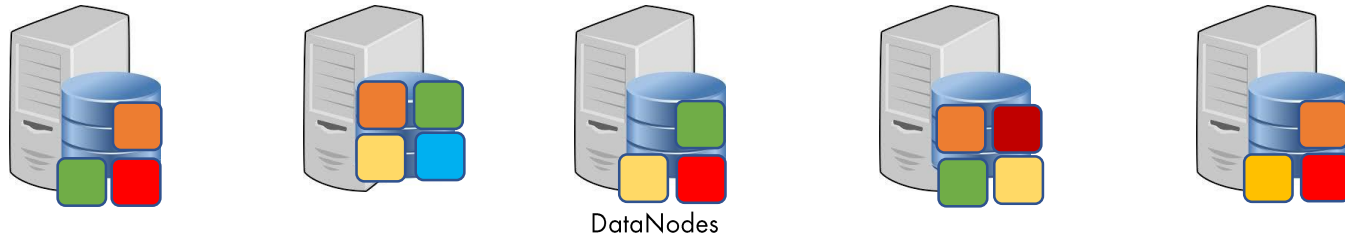
Hadoop Software Stack

HDFS Architecture

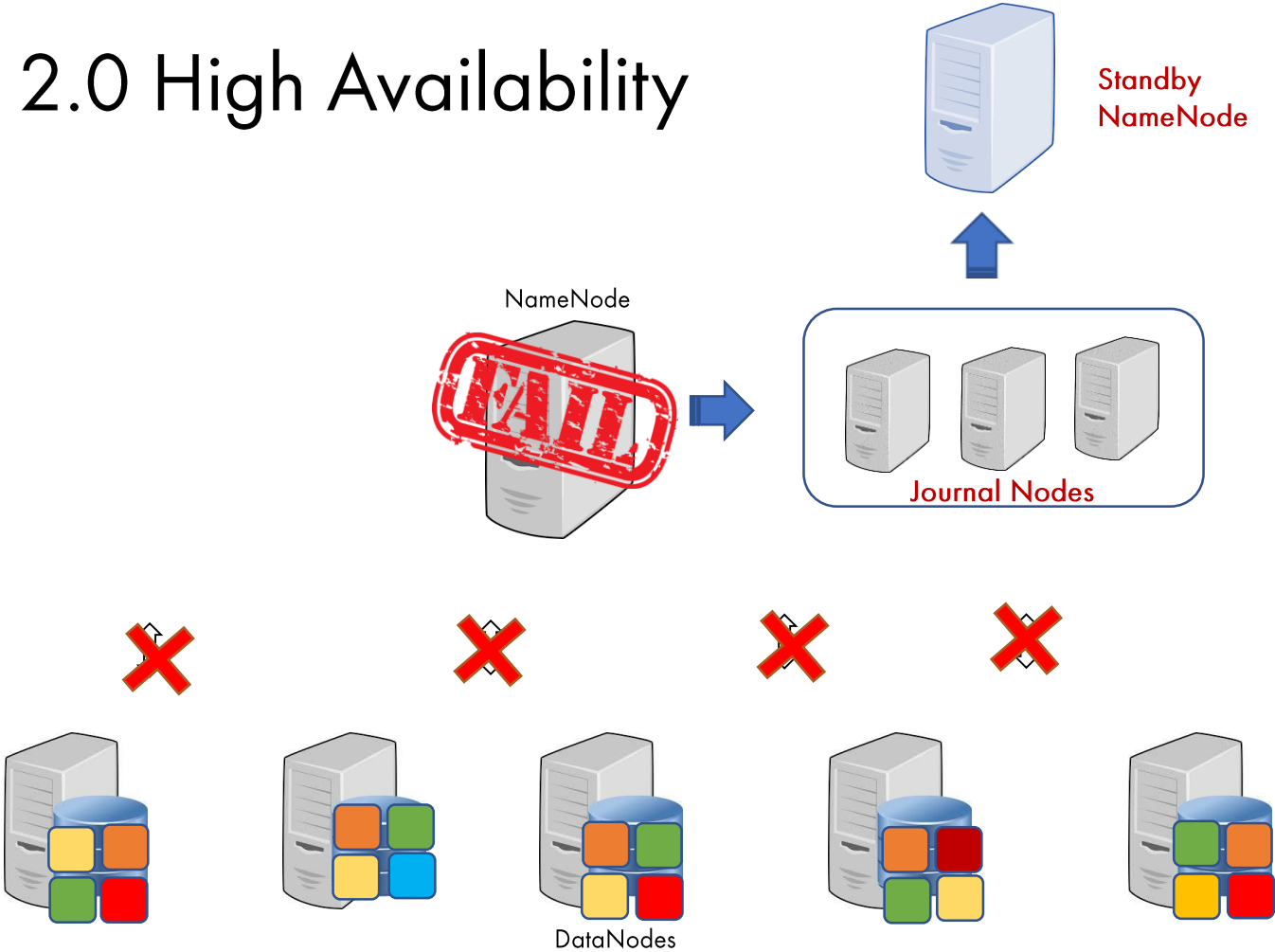


File System Metadata

File	Blocks Mappings
File 1	Blk 1 → DN1, Blk 2 → DN5, Blk 3 → DN3
File 2	Blk 1 → DN1, Blk 2 → DN4
File 3	Blk 1 → DN1, Blk 2 → DN2, Blk 3 → DN3
File 4	Blk 1 → DN100
File 5	Blk 1 → DN4, Blk 2 → DN2, Blk 3 → DN9
.....	
FileN	Blk 1 → DN2, Blk 2 → DN8



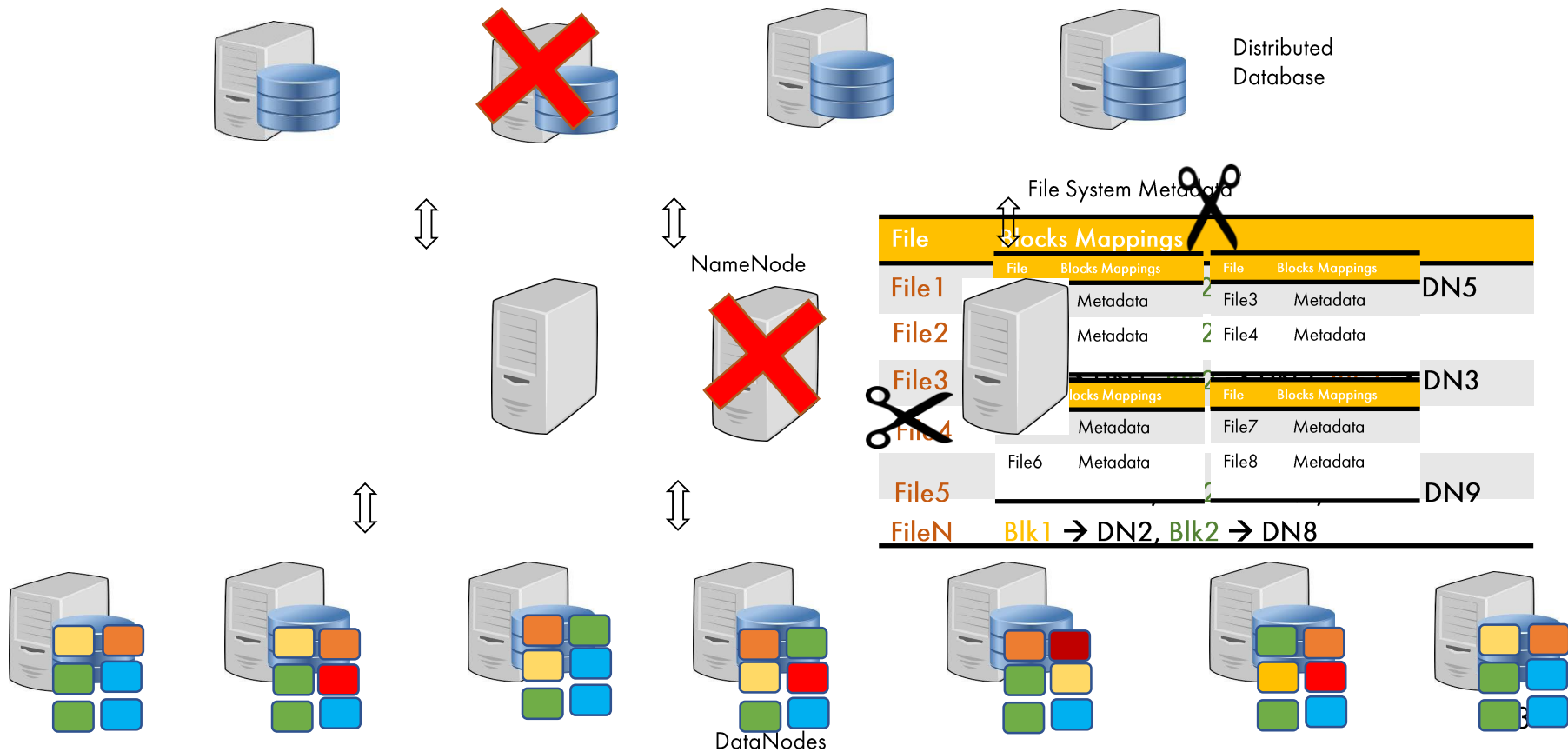
HDFS 2.0 High Availability



HDFS Limitations

- HDFS has been scaled to store **100 PB – 200 PB** on 4000 – 5000 datanodes
- Namespace size upper bound: **~ 500 million files**
- At most **70 – 80 thousand** file system operations / sec

HopsFS Architecture



NewSQL DB

MySQL Cluster: Network Database Engine (NDB)

- Open Source
- Commodity Hardware
 - Scales to 48 database nodes
 - 200 Million Read Ops/Sec* using NDB native API
 - Read Committed Transaction Isolation
 - Row-level Locking
 - User-defined partitioning

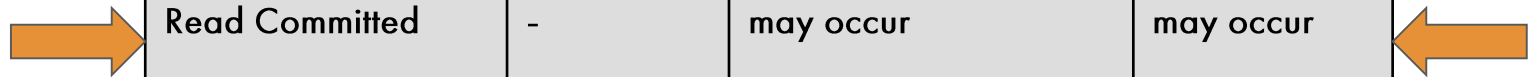
*<https://www.mysql.com/why-mysql/benchmarks/mysql-cluster/>

Transaction Isolation

Transaction Isolation Levels

MySQL Cluster Network Database Engine only supports Read-Committed Transaction Isolation Level

Isolation level	Dirty reads	Non-repeatable reads	Phantoms
Read Uncommitted	may occur	may occur	may occur
Read Committed	-	may occur	may occur
Repeatable Read	-	-	may occur
Serializable	-	-	-



Read Committed Transaction Isolation Level

Transaction 1

Start Tx
Read x (x = 1)

Read x (x = 1)

Read x (x = 10)

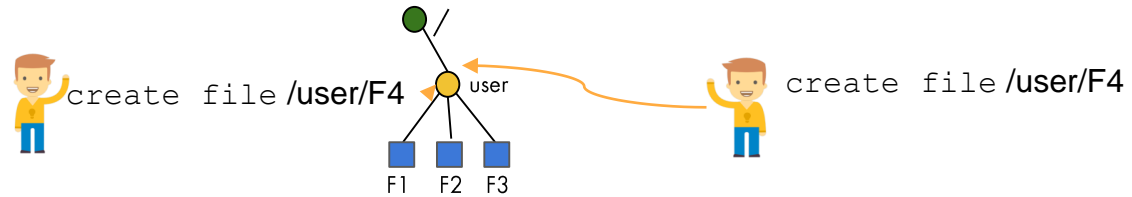
**The value of X has changed
before the transaction is
committed**

Transaction 2

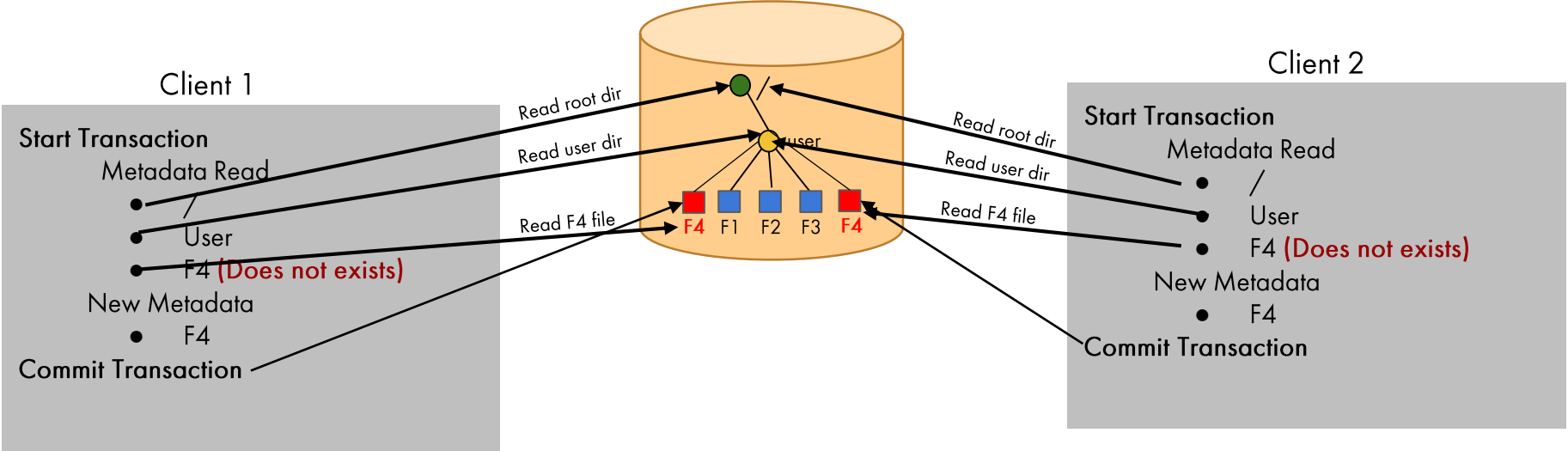
Start Tx
Read x (x = 1)
Update x = 10

Commit

Read Committed Transaction Isolation



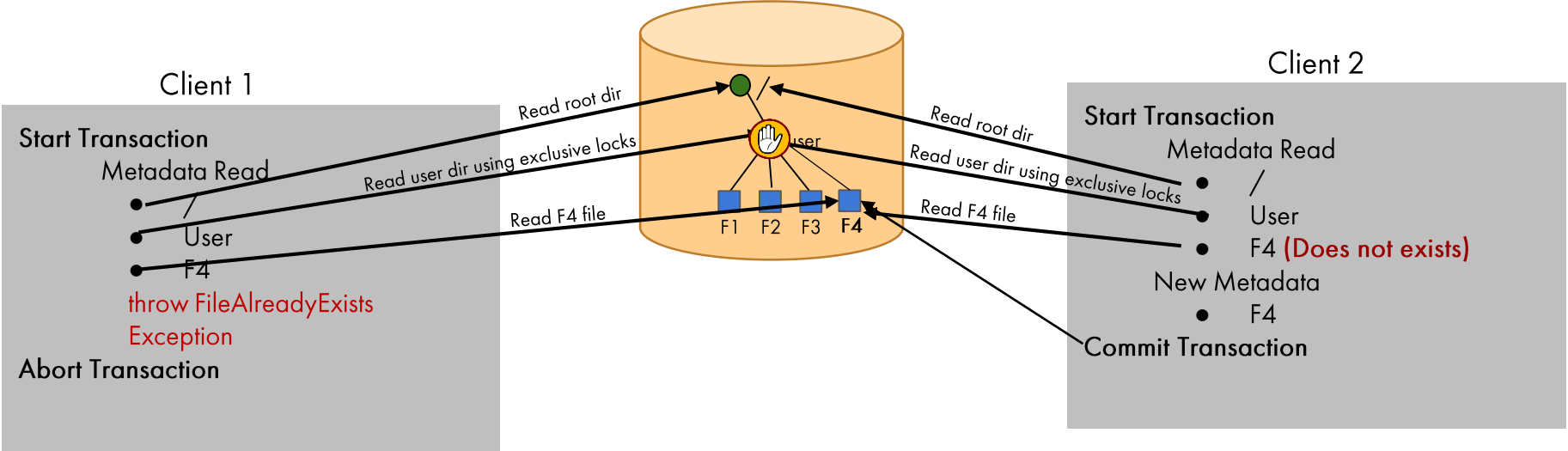
Read Committed Transaction Isolation



Read Committed Transaction Isolation

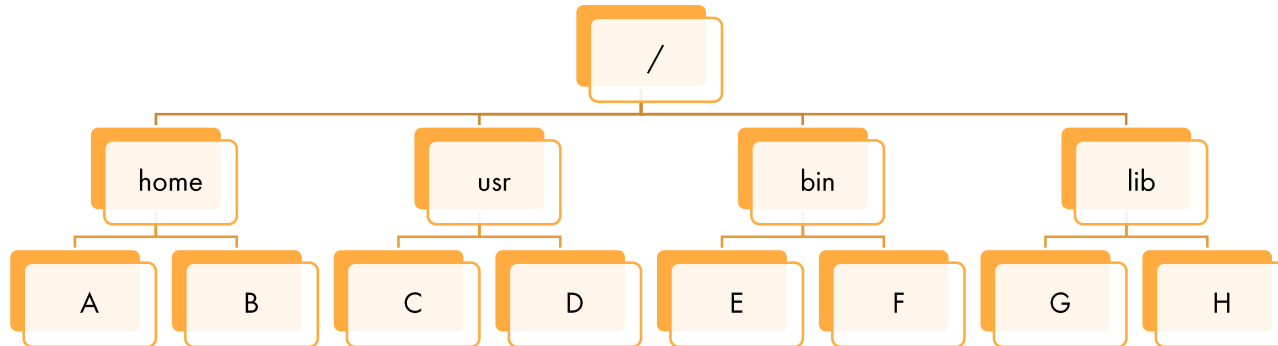
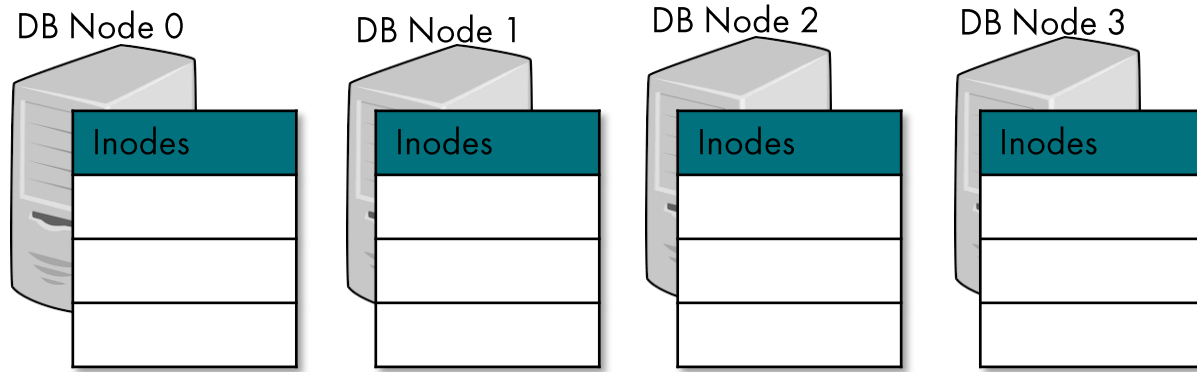
Use row level locking to serialize conflicting file operation

Read Committed Transaction Isolation With Locking



Database Operations & Data Partitioning

Distributed Metadata Design



Distributed Metadata Design

TC		
ID	name	PID
1	home	0
7	C	2
10	F	3
...

DB Node 0

TC		
ID	name	PID
5	A	1
8	D	2
4	lib	0
...

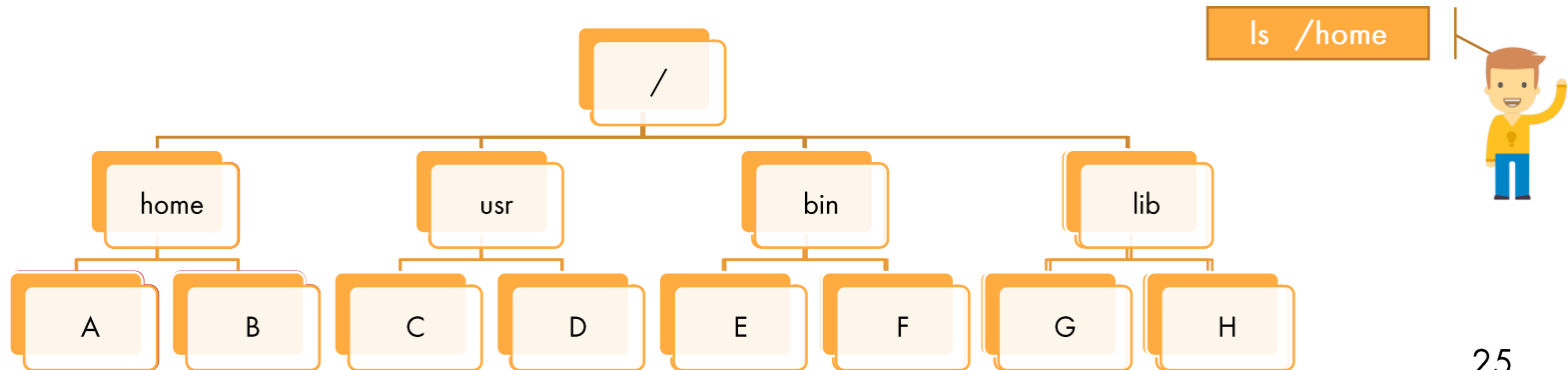
DB Node 1

TC		
ID	name	PID
6	B	1
3	bin	0
11	G	4
...

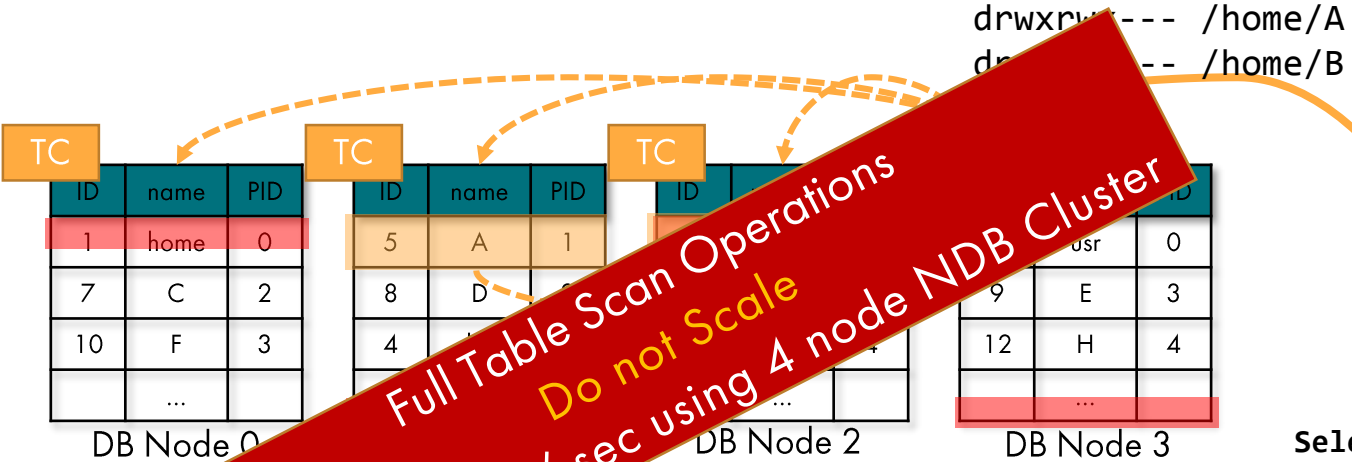
DB Node 2

TC		
ID	name	PID
2	usr	0
9	E	3
12	H	4
...

DB Node 3



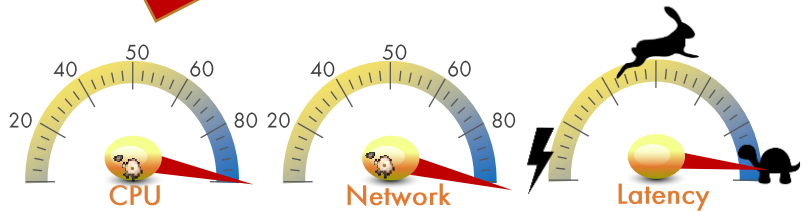
Distributed Full Table Scans



drwxrwxr-x /home/A
drwxrwxr-x /home/B

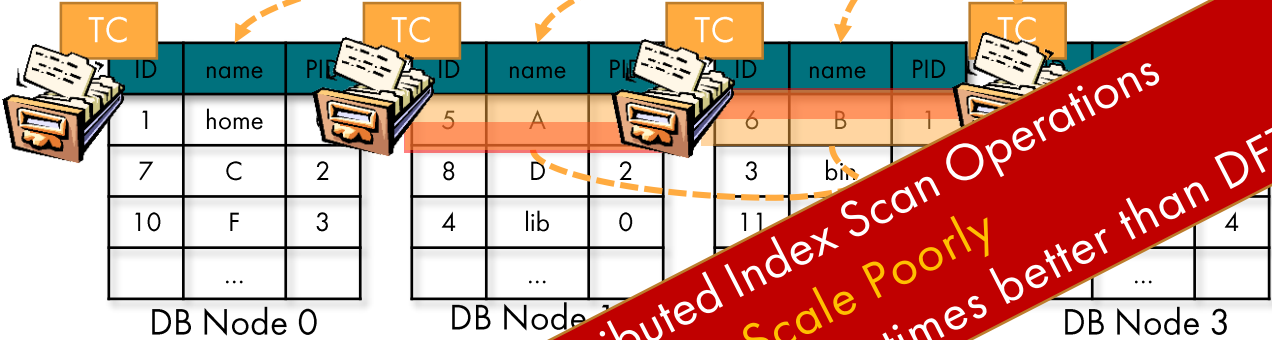
Select * from Inodes where PID = 1

ls /home



Distributed Index Scan Operations

drwxrwx--- /home/A
 drwxrwx--- /home



Distributed Index Scan Operations
 Scale Poorly
 210 K ops/sec (590 times better than DFTS)

```
Select * from Inodes where
PID = 1
WITH(INDEX(...))
```

ls -l /home



HopsFS

- **Uses NewSQL Relational Database that allows**
 1. **User Defined Partitioning (UDP):** Take control of how the data is distributed across different database nodes
 2. **Distribution Aware Transactions (DAT):** Take control over which Transaction Coordinator handles which file system operations.

Solution (User Defined Partitioning)

DB Node 0

ID	Name	PID
1	/	0
9	E	4
10	F	4
15	K	8
15	L	8

DB Node 1

ID	Name	PID
2	home	1
3	usr	1
4	bin	1
11	G	5
12	H	5

DB Node 2

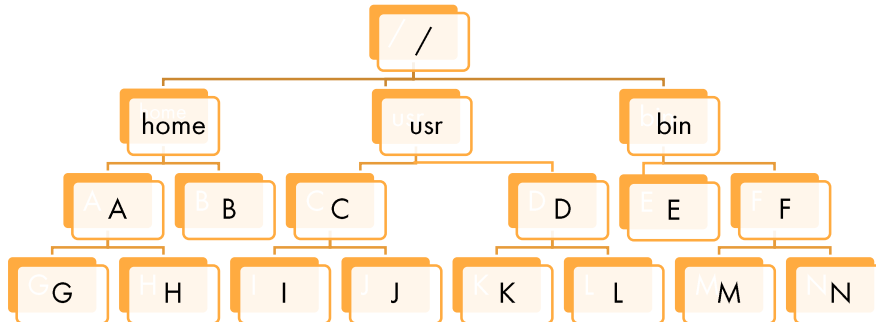
ID	Name	PID
5	A	2
6	B	2
16	M	10
17	N	10

DB Node 3

ID	Name	PID
7	C	3
8	D	3
13	I	7
14	J	7

Hash Fn

PID % 4 = Partition No



Solution (Partition Pruned Index Scan Operations)

```
drwxrwx--- /home/A
drwxrwx--- /home/B
```

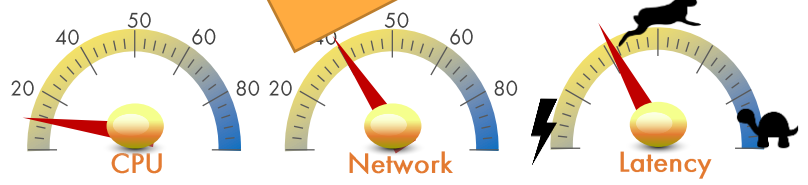
TC	Name	PID	TC	Name	PID	TC	Name	PID
1	/	0	2	home	1	5		
9	E	4	3	usr	1			
10	F	4	4	bin				7
15	K	8	11					
15	L	8						

DB Node 1 DB Node 2 DB Node 3

Partition Pruned Index Scan Scales Well

Select * from Inodes where
PID = 2
 WITH(INDEX(...))

ls -l /home



Solution (Distribution Aware Transactions)

```
drwxrwx--- /home/A
drwxrwx--- /home/B
```

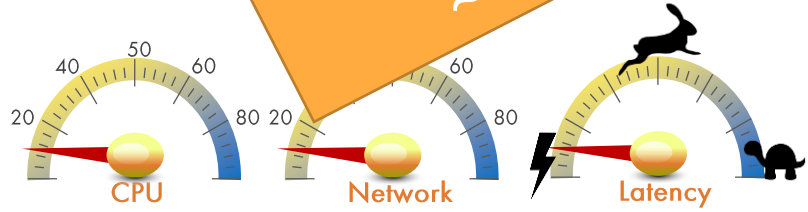
TC			TC			TC			TC		
ID	Name	PID	ID	Name	PID	ID	Name	PID	ID	Name	PID
1	/	4	2	home	1	5	A	1	6	B	1
9	E	4	3	usr	1	7	C	1	8	D	1
10	F	4	4	bin	1	9	E	1	10	F	1
15	K	8	11	G	1	11	G	1	12	H	1
15	L	8	12	H	1	13	I	1	14	J	1

DB Node 0 DB Node 3

Distribution Aware Transactions &
 Partition Pruned Index Scan
Scales Very Well
 ~ 1 Million ops/sec (5X DIS)

Start Transaction on Node 2
 Select * from Inodes where
 PID = 2
 WITH(INDEX(name))

ls -l /home



Transactional FS Operations

- File System Operation
 - Distributed Transaction
 - START DISTRIBUTION AWARE TRANSACTION
 - Primary Key Ops
 - Partition Pruned Index Scan Ops
 - Batching and Caching
 - COMMIT TRANSACTION

Solution (Contd.)

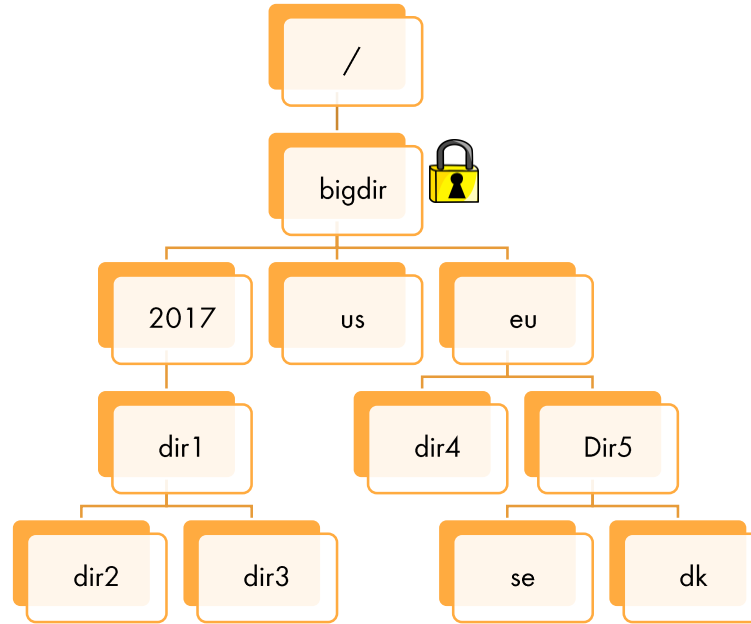
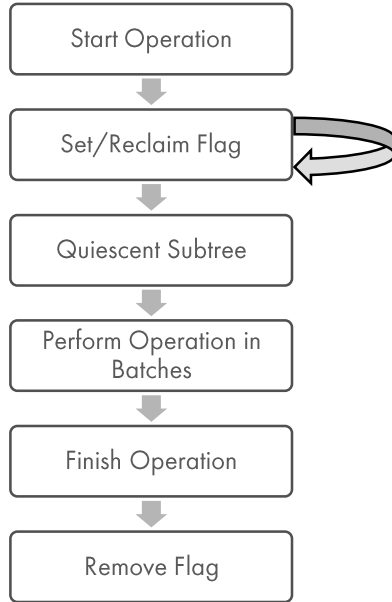
- ~40 Tables & hundreds of FS operations.
- Most NewSQL databases do not provide serializable transaction isolation level
 - HopsFS uses **read committed** transaction isolation level and **row level locks**
- Avoiding deadlocks in file system operation.
 - **Total Order Locking**
- Implementing large file system operations that do not fit in a transaction



Large File System Operations

Subtree Operations

Subtree operation Stages



Failures during Subtree Operations

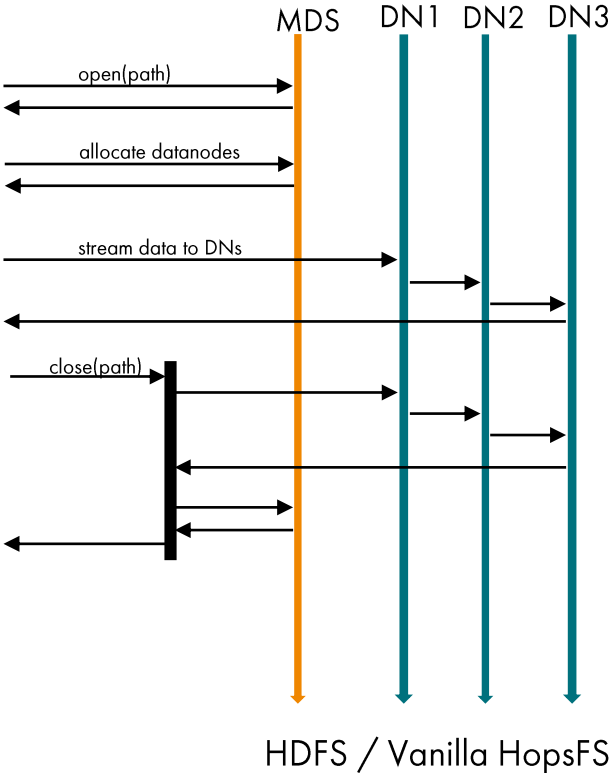
All subtree operations are implemented in such a way that if the operations fail halfway, then the namespace is not left in an inconsistent state.

HopsFS Performance

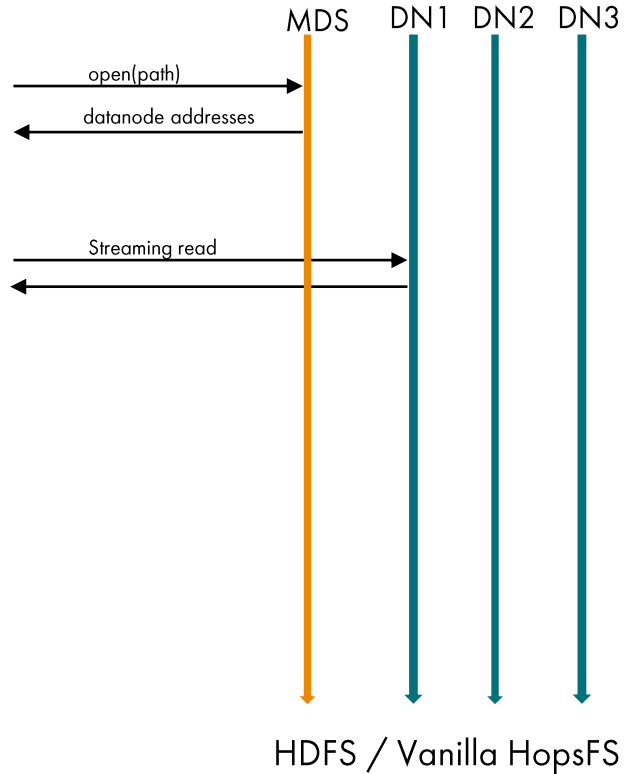
- **Throughput**
 - **16X** the throughput of HDFS (Spotify Workload).
 - **38X** the throughput of HDFS for 20% write intensive workload
- **Low Latency**
 - Identical avg op latency (~3ms) for small number (50) of clients
 - **10X** lower latency for large number (6500) of clients
- **Metadata Scaling**
 - **37X** more metadata.

Support for Small Files

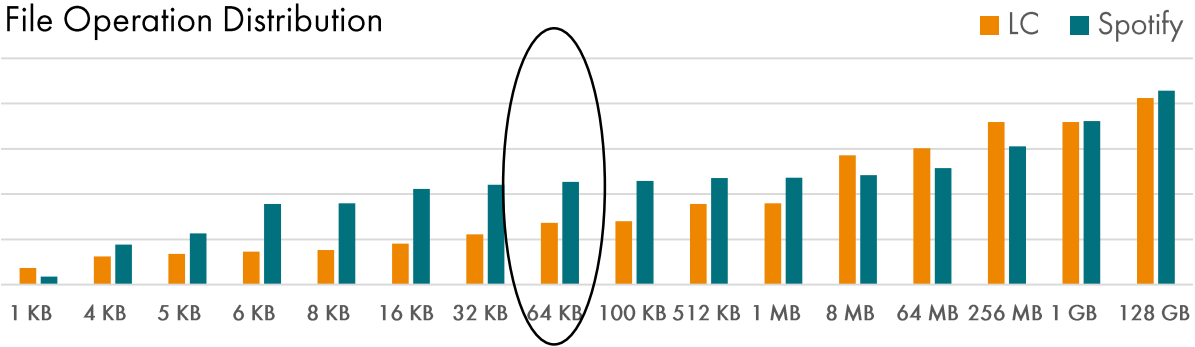
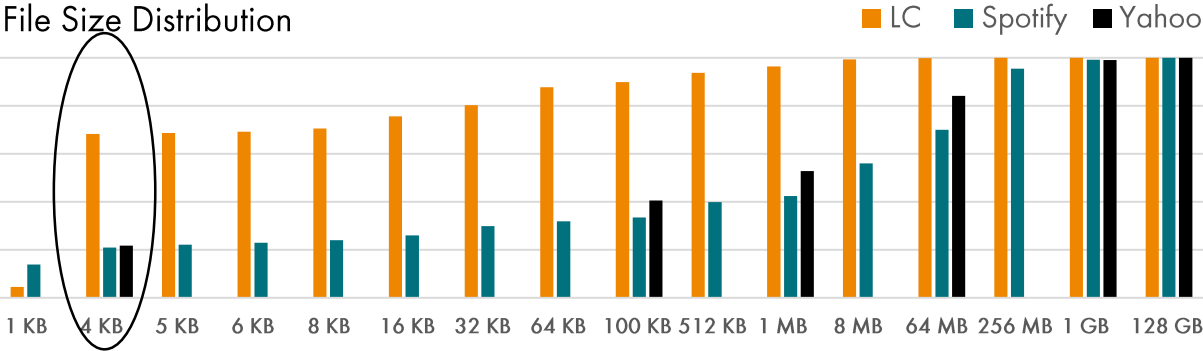
HDFS/HopsFS Write Operation



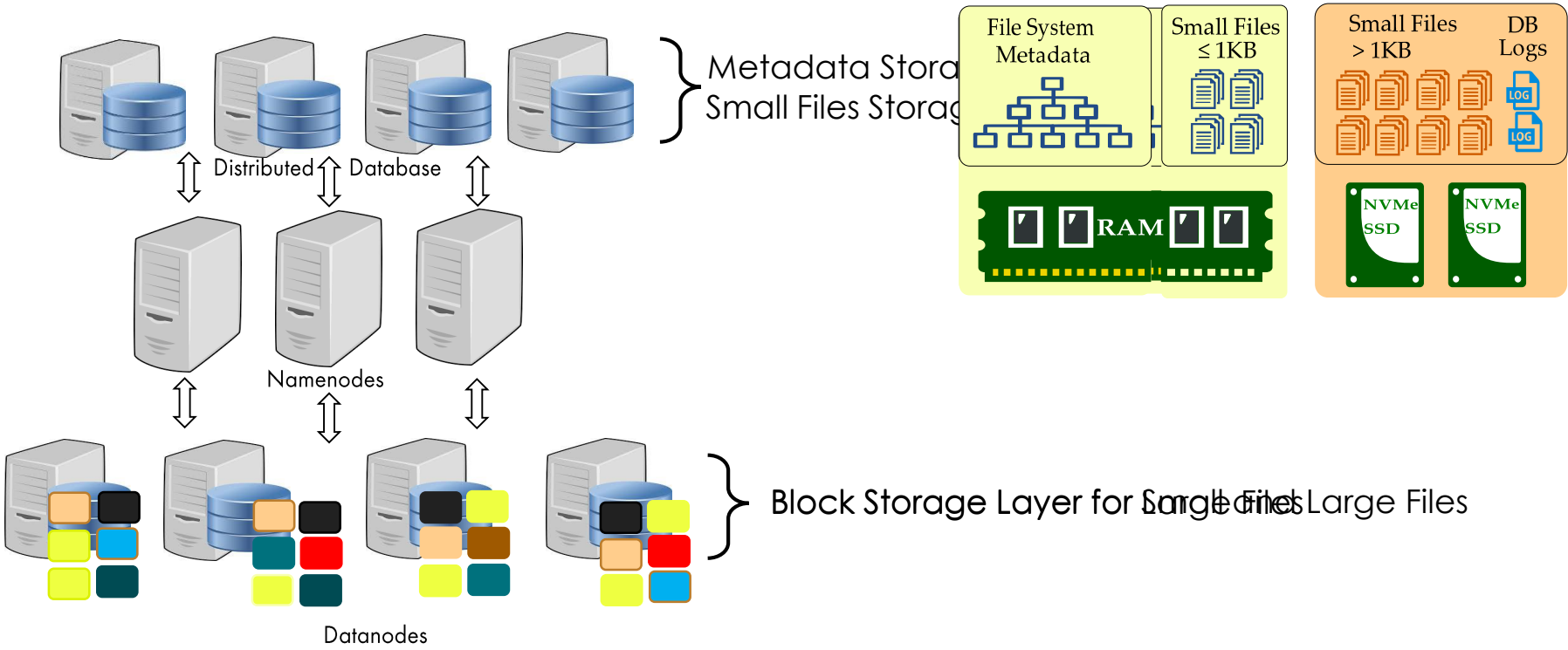
HDFS/HopsFS Read Operation



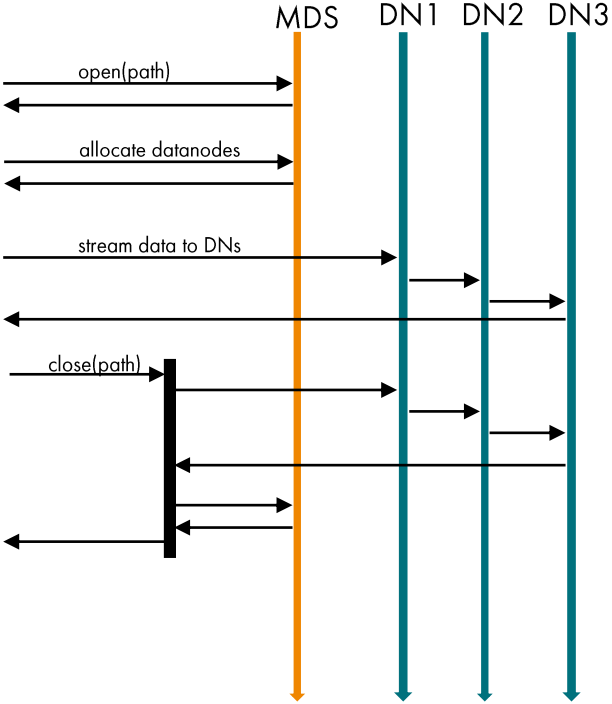
Prevalence of Small Files In Hadoop



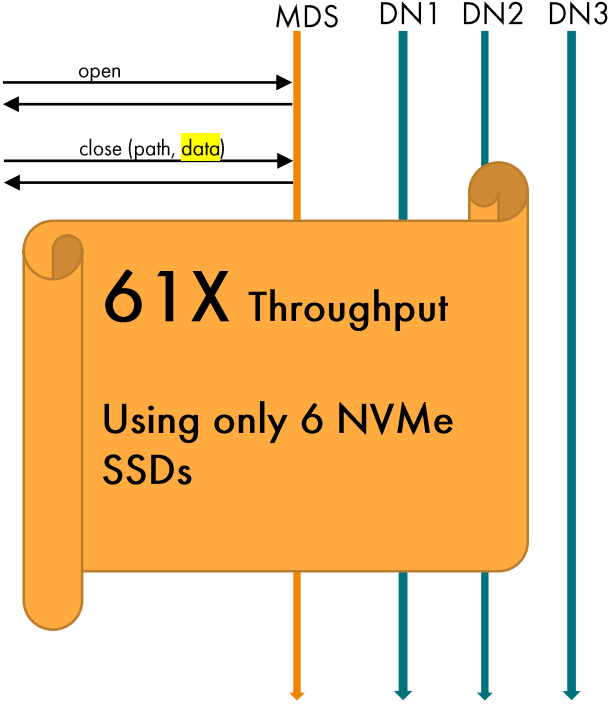
Small Files' Support in HopsFS



Optimizing Write Operation For Small Files

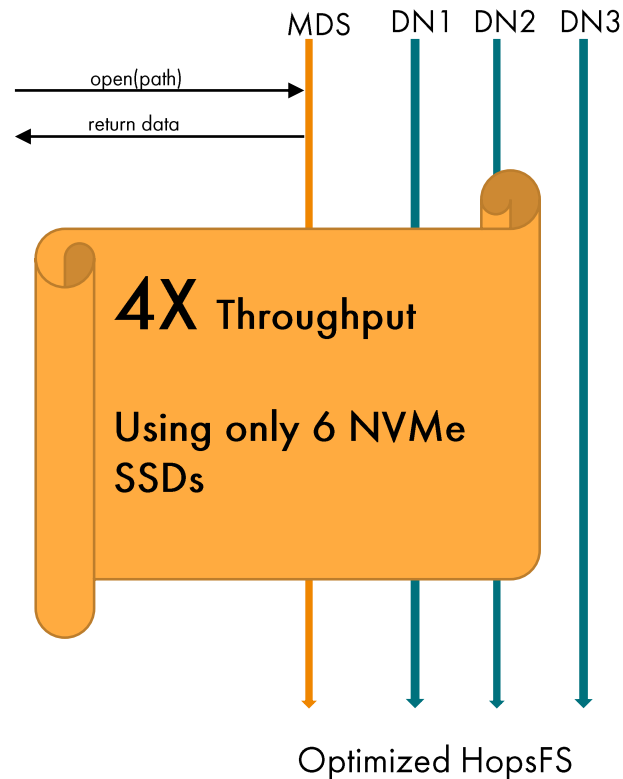
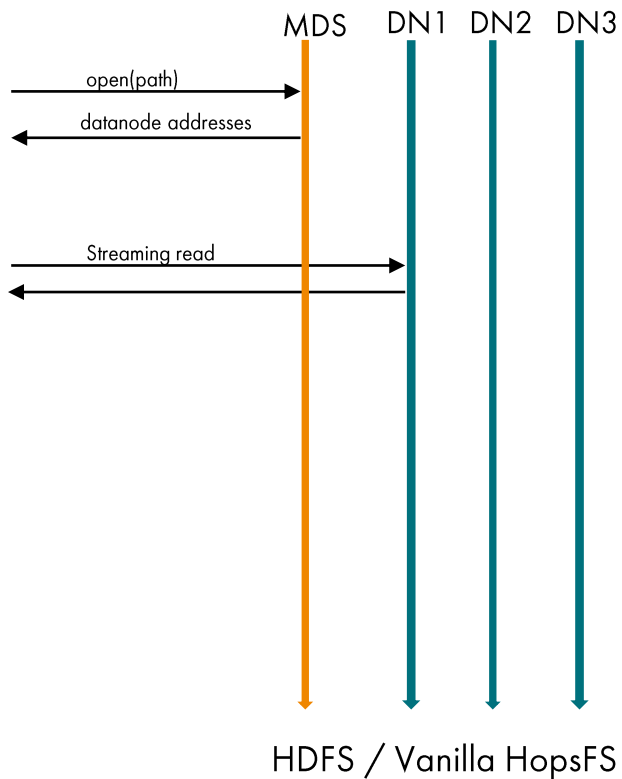


HDFS / Vanilla HopsFS



Optimized HopsFS

Optimizing Read Operations For Small Files



Questions



<http://www.hops.io>



<http://github.com/hopshadoop>



@hopshadoop

Read More

Scaling Distributed Hierarchical File Systems Using NewSQL Databases
Salman Niazi. Ph.D. Thesis. KTH Royal Institute of Technology

Scaling hierarchical file system metadata using newsql databases
S Niazi, M Ismail, S Haridi, J Dowling, S Grohsschmiedt, M Ronström
15th USENIX Conference on File and Storage Technologies (FAST 17), 89-104

Scaling HDFS to more than 1 million operations per second with HopsFS
M Ismail, S Niazi, M Ronström, S Haridi, J Dowling
2017 17th IEEE/ACM International Symposium on Cluster, Cloud and Grid ...

Size Matters: Improving the Performance of Small Files in Hadoop
S Niazi, M Ronström, S Haridi, J Dowling
Proceedings of the 19th International Middleware Conference, 26-39