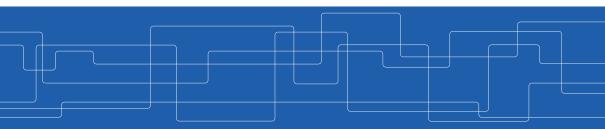


#### Introduction to Data Stream Processing

Amir H. Payberah payberah@kth.se 2022-09-22





#### The Course Web Page

# https://id2221kth.github.io

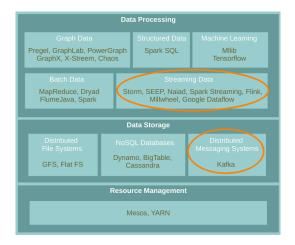


The Questions-Answers Page

# https://tinyurl.com/bdenpwc5



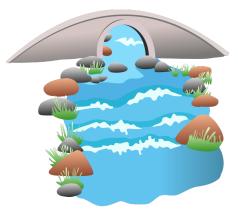
#### Where Are We?





### Stream Processing (1/3)

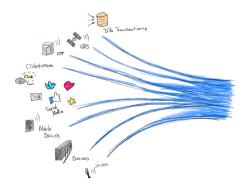
Stream processing is the act of continuously incorporating new data to compute a result.





### Stream Processing (2/3)

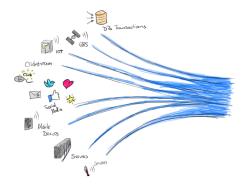
- ► The input data is unbounded.
  - A series of events, no predetermined beginning or end.





### Stream Processing (2/3)

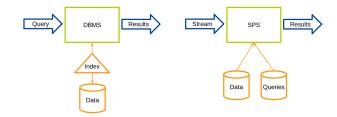
- The input data is unbounded.
  - A series of events, no predetermined beginning or end.
  - E.g., credit card transactions, clicks on a website, or sensor readings from IoT devices.





### Stream Processing (3/3)

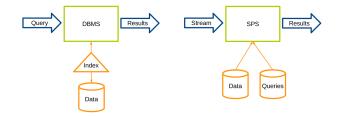
- Database Management Systems (DBMS): data-at-rest analytics
  - Store and index data before processing it.
  - Process data only when explicitly asked by the users.





### Stream Processing (3/3)

- Database Management Systems (DBMS): data-at-rest analytics
  - Store and index data before processing it.
  - Process data only when explicitly asked by the users.
- ► Stream Processing Systems (SPS): data-in-motion analytics
  - Processing information as it flows, without storing them persistently.





### Stream Processing Systems Stack

Pr	U	ີ	C	3	3		u

Spark Streaming, Flink, Storm, Google Dataflow

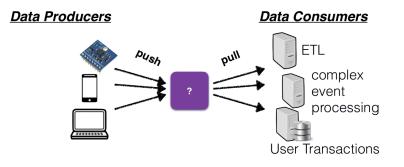
Storage					
Partitioned Logs	Messaging Systems				
Apache Kafka, Amazon Kinesis Twitter Distributed Log	Google Cloud Pub/Sub, RabbitMQ ActiveMQ, Azure Service Bus				



# Data Stream Storage

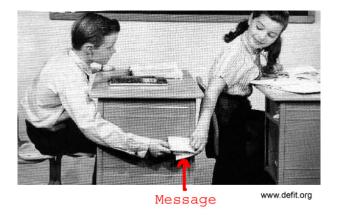


► We need disseminate streams of events from various producers to various consumers.





#### Messaging systems





#### What is Messaging System?

• Messaging system is an approach to notify consumers about new events.

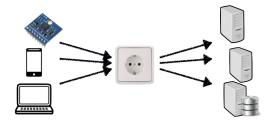


### What is Messaging System?

- ▶ Messaging system is an approach to notify consumers about new events.
- Messaging systems
  - Direct messaging
  - Message brokers

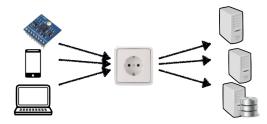


- ▶ Necessary in latency critical applications (e.g., remote surgery).
- A producer sends a message containing the event, which is **pushed** to consumers.



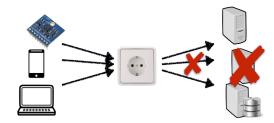


- ▶ Necessary in latency critical applications (e.g., remote surgery).
- A producer sends a message containing the event, which is **pushed** to consumers.
- Both consumers and producers have to be online at the same time.



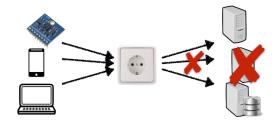


▶ What happens if a consumer crashes or temporarily goes offline? (not durable)



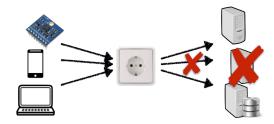


- ► What happens if a consumer crashes or temporarily goes offline? (not durable)
- ▶ What happens if producers send messages faster than the consumers can process?



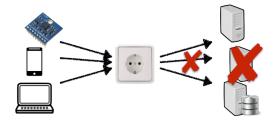


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  - Dropping messages
  - Backpressure





- ► What happens if a consumer crashes or temporarily goes offline? (not durable)
- ▶ What happens if producers send messages faster than the consumers can process?
  - Dropping messages
  - Backpressure
- ▶ We need message brokers that can log events to process at a later time.



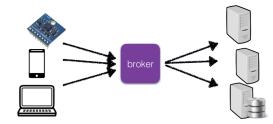




[https://bluesyemre.com/2018/10/16/thousands-of-scientists-publish-a-paper-every-five-days]

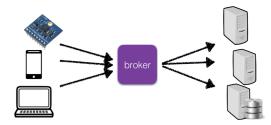


- A message broker decouples the producer-consumer interaction.
- ▶ It runs as a server, with producers and consumers connecting to it as clients.



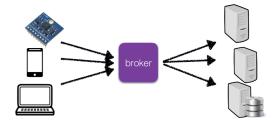


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- Producers write messages to the broker, and consumers receive them by reading them from the broker.





- ► A message broker decouples the producer-consumer interaction.
- ► It runs as a server, with producers and consumers connecting to it as clients.
- Producers write messages to the broker, and consumers receive them by reading them from the broker.
- Consumers are generally asynchronous.





#### ► In typical message brokers, once a message is consumed, it is deleted.



Partitioned Logs

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- ► Log-based message brokers durably store all events in a sequential log.



#### Partitioned Logs

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- ► Log-based message brokers durably store all events in a sequential log.
- A log is an append-only sequence of records on disk.



#### Partitioned Logs

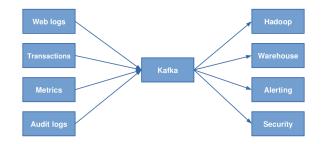
- ► In typical message brokers, once a message is consumed, it is deleted.
- ► Log-based message brokers durably store all events in a sequential log.
- A log is an append-only sequence of records on disk.
- ► A producer sends a message by appending it to the end of the log.
- ► A consumer receives messages by reading the log sequentially.



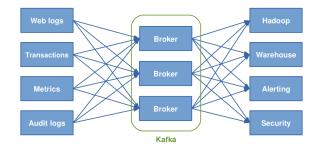
# Kafka - A Log-Based Message Broker



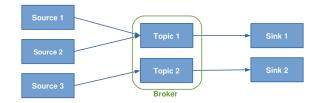




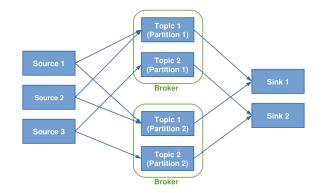




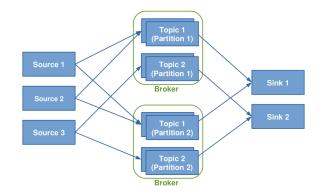














### Logs, Topics and Partition (1/6)

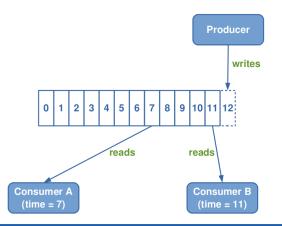
- ► Kafka is about logs.
- Topics are queues: a stream of messages of a particular type

jkreps-mn:~ jkreps\$ tail -f -n 20 /v	
::1 [23/Mar/2014:15:07:00 -0700]	"GET /images/apache_feather.gif HTTP/1.1" 200 4128
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/producer_consumer.png HTTP/1.1" 200 86
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/log_anatomy.png HTTP/1.1" 200 19579
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/consumer-groups.png HTTP/1.1" 200 2682
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/log_compaction.png HTTP/1.1" 200 41414
::1 [23/Mar/2014:15:07:04 -0700]	"GET /documentation.html HTTP/1.1" 200 189893
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/log_cleaner_anatomy.png HTTP/1.1" 200
::1 [23/Mar/2014:15:07:04 -0700]	"GET /images/mirror-maker.png HTTP/1.1" 200 17054
::1 [23/Mar/2014:15:08:07 -0700]	"GET /documentation.html HTTP/1.1" 200 189937
::1 [23/Mar/2014:15:08:07 -0700]	
::1 [23/Mar/2014:15:08:07 -0700]	"GET /images/kafka_logo.png HTTP/1.1" 304 -
::1 [23/Mar/2014:15:08:07 -0700]	"GET /images/producer_consumer.png HTTP/1.1" 304 -
::1 [23/Mar/2014:15:08:07 -0700]	"GET /images/log_anatomy.png HTTP/1.1" 304 -
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::1 [23/Mar/2014:15:09:55 -0700]	"GET /documentation.html HTTP/1.1" 200 195264

0 1	2 3 4	56	7 8 9	10 11 12
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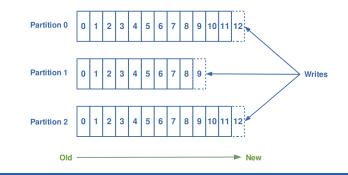
• Each message is assigned a sequential id called an offset.





## Logs, Topics and Partition (3/6)

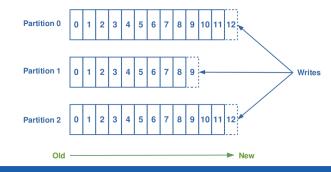
- ► Topics are logical collections of partitions (the physical files).
  - Ordered
  - Append only
  - Immutable





## Logs, Topics and Partition (4/6)

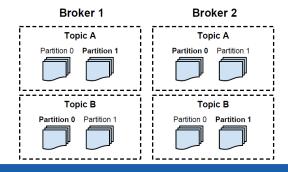
- Ordering is only guaranteed within a partition for a topic.
- Messages sent by a producer to a particular topic partition will be appended in the order they are sent.
- ▶ A consumer instance sees messages in the order they are stored in the log.





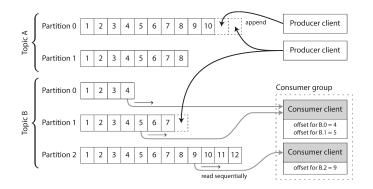
### Logs, Topics and Partition (5/6)

- ▶ Partitions of a topic are replicated: fault-tolerance
- A broker contains some of the partitions for a topic.
- One broker is the leader of a partition: all writes and reads must go to the leader.



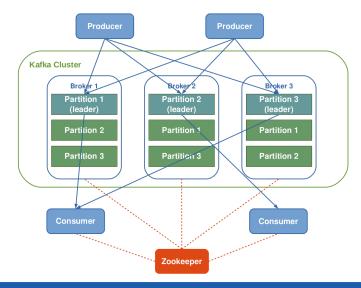


## Partitioned Logs (6/6)





#### Kafka Architecture





• Kafka uses **Zookeeper** for the following tasks:





- Kafka uses **Zookeeper** for the following tasks:
- Detecting the addition and the removal of brokers and consumers.
- Keeping track of the consumed offset of each partition.





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#### State in Kafka

- Brokers are sateless: no metadata for consumers-producers in brokers.
- Consumers are responsible for keeping track of offsets.
- ▶ Messages in queues expire based on pre-configured time periods (e.g., once a day).



Kafka guarantees that messages from a single partition are delivered to a consumer in order.



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- ► There is no guarantee on the ordering of messages coming from different partitions.
- ► Kafka only guarantees at-least-once delivery.



# Start the ZooKeeper

zookeeper-server-start.sh config/zookeeper.properties



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# Start the Kafka server
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# Produce messages and send them to the topic "avg"
kafka-console-producer.sh --topic avg --bootstrap-server localhost:9092



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# Consume the messages sent to the topic "aug"
kafka-console-consumer.sh --topic avg --from-beginning --bootstrap-server localhost:9092



# Data Stream Processing





- Data stream is unbound data, which is broken into a sequence of individual tuples.
- A data tuple is the atomic data item in a data stream.
- ► Can be structured, semi-structured, and unstructured.



#### Streaming Data Processing Design Points

- Continuous vs. micro-batch processing
- Record-at-a-Time vs. declarative APIs
- Event time vs. processing time
- Windowing



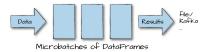
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- Micro-batch systems
  - Batch engines
  - Slicing up the unbounded data into a sets of bounded data, then process each batch.





## Streaming Data Processing Patterns

- Micro-batch systems
  - Batch engines
  - Slicing up the unbounded data into a sets of bounded data, then process each batch.



- Continuous processing-based systems
  - Each node in the system continually listens to messages from other nodes and outputs new updates to its child nodes.





#### Streaming Data Processing Design Points

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#### Record-at-a-Time vs. Declarative APIs

- Record-at-a-Time API (e.g., Storm)
  - Low-level API
  - Passes each event to the application and let it react.
  - Useful when applications need full control over the processing of data.
  - Complicated factors, such as maintaining state, are governed by the application.



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  - Useful when applications need full control over the processing of data.
  - Complicated factors, such as maintaining state, are governed by the application.
- ► Declarative API (e.g., Spark streaming, Flink, Google Dataflow)
  - Aapplications specify what to compute not how to compute it in response to each new event.



#### Streaming Data Processing Design Points

- Continuous vs. micro-batch processing
- ► Record-at-a-Time vs. declarative APIs
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#### Event Time vs. Processing Time (1/2)

- Event time: the time at which events actually occurred.
  - Timestamps inserted into each record at the source.
- ▶ Prcosseing time: the time when the record is received at the streaming application.



### Event Time vs. Processing Time (2/2)

- Ideally, event time and processing time should be equal.
- Skew between event time and processing time.



[https://www.oreilly.com/ideas/the-world-beyond-batch-streaming-101]



#### Streaming Data Processing Design Points

- Continuous vs. micro-batch processing
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### Windowing (1/2)

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  - Count-based policy: the maximum number of tuples a window buffer can hold
  - Delta-based policy: a delta threshold in a tuple attribute
  - Punctuation-based policy: a punctuation is received
  - Time-based policy: based on processing or event time period



## Windowing (2/2)

#### Two types of windows: tumbling and sliding



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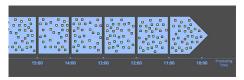
- ► Sliding window: supports incremental operations.
  - When the buffer fills up, older tuples are evicted.

1 21 321	4321 5432 6543
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### Windowing by Processing Time

- The system buffers up incoming data into windows until some amount of processing time has passed.
- ► E.g., five-minute fixed windows

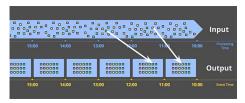


[https://www.oreilly.com/ideas/the-world-beyond-batch-streaming-101]



### Windowing by Event Time

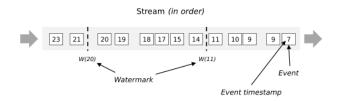
- Reflect the times at which events actually happened.
- Handling out-of-order evnets.



[https://www.oreilly.com/ideas/the-world-beyond-batch-streaming-101]

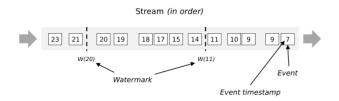


► Watermarking helps a stream processing system to deal with lateness.



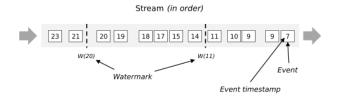


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- Watermarks flow as part of the data stream and carry a timestamp t.



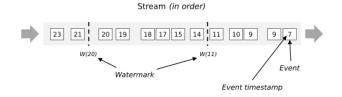


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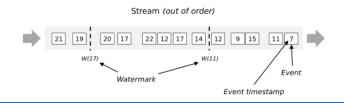


- ► Watermarking helps a stream processing system to deal with lateness.
- Watermarks flow as part of the data stream and carry a timestamp t.
- ► A watermark is a threshold to specify how long the system waits for late events.
- ► Streaming systems uses watermarks to measure progress in event time.



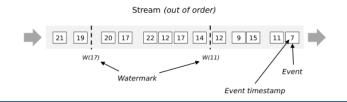


- A W(t) declares that event time has reached time t in that stream
  - There should be no more elements from the stream with a timestamp  $t^\prime \leq t.$



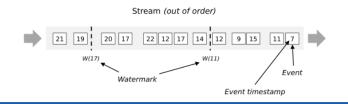


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- ▶ If an arriving event lies within the watermark, it gets used to update a query.
- Streaming programs may explicitly expect some late elements.



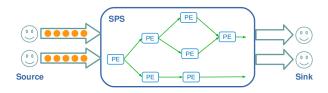


# Streaming Data Processing Model



### Streaming Data Processing

- ► The tuples are processed by the application's operators or processing element (PE).
- A PE is the basic functional unit in an application.
  - A PE processes input tuples, applies a function, and outputs tuples.
  - A set of PEs and stream connections, organized into a data flow graph.





### PEs States (1/3)

- ► A PE can either maintain internal state across tuples while processing them, or process tuples independently of each other.
- Stateful vs. stateless tasks



## PEs States (2/3)

Stateless tasks: do not maintain state and process each tuple independently of prior history, or even from the order of arrival of tuples.



### PEs States (2/3)

- Stateless tasks: do not maintain state and process each tuple independently of prior history, or even from the order of arrival of tuples.
- Easily parallelized.
- No synchronization.
- ▶ Restart upon failures without the need of any recovery procedure.



## PEs States (3/3)

 Stateful tasks: involves maintaining information across different tuples to detect complex patterns.



### PEs States (3/3)

- Stateful tasks: involves maintaining information across different tuples to detect complex patterns.
- A PE is usually a synopsis of the tuples received so far.
- A subset of recent tuples kept in a window buffer.



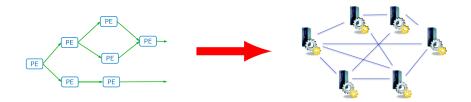
### Job and Job Management

- ► At runtime, an application is represented by one or more jobs.
- ► Jobs are deployed as a collection of PEs.
- ► Job management component must identify and track individual PEs, the jobs they belong to, and associate them with the user that instantiated them.



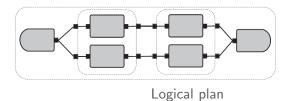
### Logical Plan vs. Physical Plan (1/2)

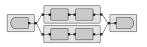
- Logical plan: a data flow graph, where the vertices correspond to PEs, and the edges to stream connections.
- Physical plan: a data flow graph, where the vertices correspond to OS processes, and the edges to transport connections.

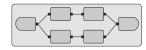


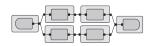


### Logical Plan vs. Physical Plan (2/2)









Different physical plans

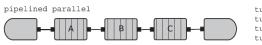


Parallelization

- ▶ How to scale with increasing the number queries and the rate of incoming events?
- ► Three forms of parallelisms.
  - Pipelined parallelism
  - Task parallelism
  - Data parallelism



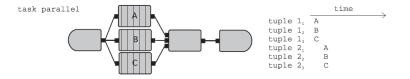
#### ► Sequential stages of a computation execute concurrently for different data items.



		time							~	
tuple	1	A	В	С					_	
tuple	2		А	В	С					
tuple	3			А	В	С				
tuple	4				А	В	С			
tuple	5					А	В	С		
tuple	6						А	В	С	

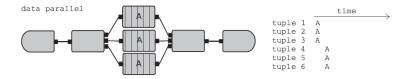


Independent processing stages of a larger computation are executed concurrently on the same or distinct data items.



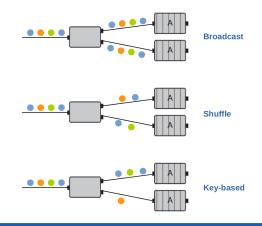


#### ► The same computation takes place concurrently on different data items.





▶ How to allocate data items to each computation instance?





# Summary





- Messaging system and partitioned logs
- Decoupling producers and consumers
- ► Kafka: distributed, topic oriented, partitioned, replicated log service
- Logs, topcs, partition
- ► Kafka architecture: producer, consumer, broker, coordinator



- ► SPS vs. DBMS
- Data stream, unbounded data, tuples
- ▶ Event-time vs. processing time
- Micro-batch vs. continues processing (windowing)
- PEs and dataflow
- Stateless vs. Stateful PEs



- J. Kreps et al., "Kafka: A distributed messaging system for log processing", NetDB 2011
- ▶ M. Zaharia et al., "Spark: The Definitive Guide", O'Reilly Media, 2018 Chapter 20
- M. Fragkoulis et al., "A Survey on the Evolution of Stream Processing Systems", 2020
- T. Akidau, "The world beyond batch: Streaming 101", https://www.oreilly.com/ideas/the-world-beyond-batch-streaming-101



# Questions?