



# NoSQL Databases

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## The Course Web Page

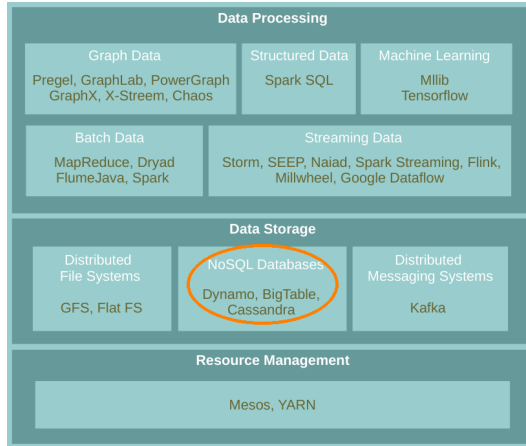
`https://id2221kth.github.io`



## The Questions-Answers Page

<https://tinyurl.com/hk7hzpw5>

# Where Are We?



# Database and Database Management System

- ▶ **Database:** an **organized** collection of **data**.
- ▶ **Database Management System (DBMS):** a **software** to **capture** and **analyze** data.





# SQL vs. NoSQL Databases

# Relational SQL Databases

- ▶ The **dominant** technology for storing **structured** data in web and business applications.
- ▶ SQL is good
  - **Rich** language and toolset
  - **Easy** to use and integrate
  - Many **vendors**
- ▶ They promise: **ACID**





# ACID Properties

## ▶ Atomicity

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- A database is in a **consistent** state before and after a transaction.

## ▶ Isolation

- Transactions can not see **uncommitted changes** in the database.

## ▶ Durability

- Changes are written to a **disk** before a database commits a transaction so that committed data cannot be lost through a power **failure**.

# SQL Databases Challenges

- ▶ **Web-based applications** caused spikes.
  - Internet-scale data size
  - High read-write rates
  - Frequent schema changes



# SQL Databases Challenges

- ▶ Web-based applications caused spikes.
  - Internet-scale data size
  - High read-write rates
  - Frequent schema changes
- ▶ RDBMS were **not** designed to be **distributed**.



▶ **Avoids:**

- Overhead of **ACID** properties
- **Complexity** of SQL query

▶ **Provides:**

- **Scalability**
- Easy and frequent **changes** to DB
- **Large** data volumes



# Availability vs. Consistency



## Availability

- ▶ **Replicating** data to improve the **availability** of data.



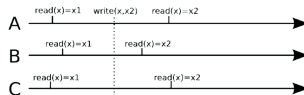
# Availability

- ▶ **Replicating** data to improve the **availability** of data.
- ▶ **Data replication**
  - Storing data in **more than one** site or node



▶ **Strong** consistency

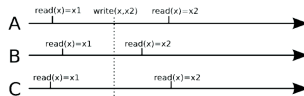
- After an update completes, any subsequent access will return the **updated value**.



# Consistency

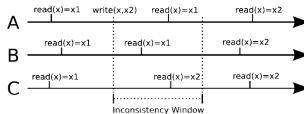
## ▶ Strong consistency

- After an update completes, any subsequent access will return the **updated value**.



## ▶ Eventual consistency

- Does **not guarantee** that subsequent accesses will return the **updated value**.
- **Inconsistency window**.
- If no new updates are made to the object, **eventually** all accesses will return the last updated value.





## Availability vs. Consistency

- ▶ The large-scale applications have to be **reliable**: consistency, availability, partition tolerance



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- ▶ Achieving **ACID** properties on large-scale applications is **challenging**.



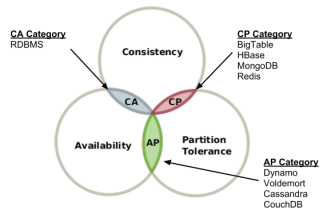
## Availability vs. Consistency

- ▶ The large-scale applications have to be **reliable**: consistency, availability, partition tolerance
- ▶ Achieving **ACID** properties on large-scale applications is **challenging**.
- ▶ **CAP** theorem

# CAP Theorem

## ► Consistency

- Consistent state of data after the execution of an operation.



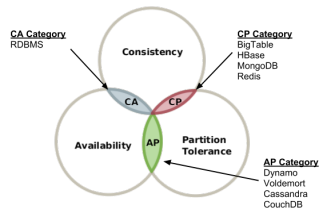
# CAP Theorem

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## ▶ Availability

- Clients can always read and write data.





# CAP Theorem

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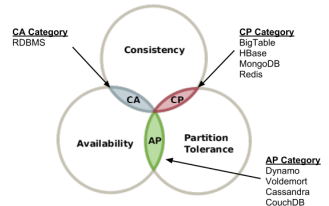
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- Continue the operation in the presence of network partitions.



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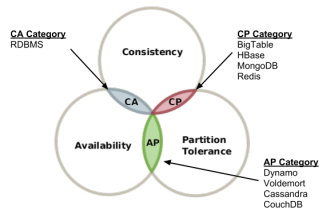
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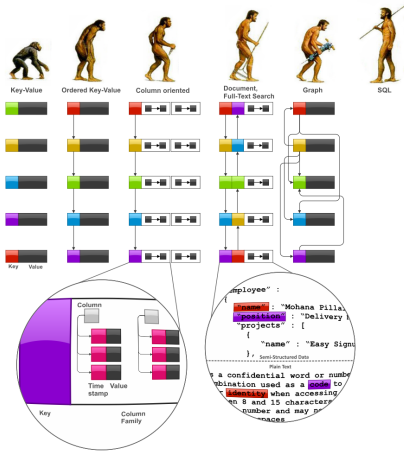
## ▶ You can choose only two!





# NoSQL Data Models

# NoSQL Data Models



[<http://highlyscalable.wordpress.com/2012/03/01/nosql-data-modeling-techniques>]



# Key-Value Data Model

- ▶ Collection of **key/value** pairs.
- ▶ **Ordered** Key-Value: processing over **key ranges**.
- ▶ **Dynamo, Scalaris, Voldemort, Riak, ...**

# Column-Oriented Data Model

- ▶ Similar to a **key/value** store, but the **value** can have multiple **attributes** (Columns).
- ▶ **Column**: a set of data **values** of a particular **type**.
- ▶ Store and process data by **column** instead of **row**.
- ▶ **BigTable**, **Hbase**, **Cassandra**, ...





# Document Data Model

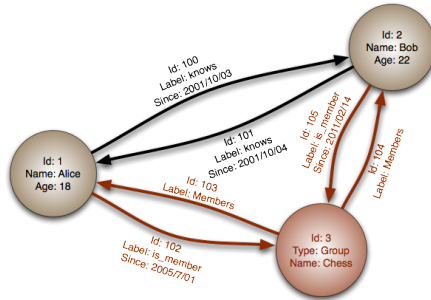
- ▶ Similar to a **column-oriented** store, but values can have **complex documents**.
- ▶ Flexible schema (XML, YAML, JSON, and BSON).
- ▶ **CouchDB, MongoDB, ...**

```
{
  FirstName: "Bob",
  Address: "5 Oak St.",
  Hobby: "sailing"
}

{
  FirstName: "Jonathan",
  Address: "15 Wanamassa Point Road",
  Children: [
    {Name: "Michael", Age: 10},
    {Name: "Jennifer", Age: 8},
  ]
}
```

# Graph Data Model

- ▶ Uses **graph** structures with **nodes**, **edges**, and **properties** to represent and store data.
- ▶ Neo4J, InfoGrid, ...



[[http://en.wikipedia.org/wiki/Graph\\_database](http://en.wikipedia.org/wiki/Graph_database)]





# BigTable

# BigTable

- ▶ Lots of (semi-)structured data at Google.
  - URLs, per-user data, geographical locations, ...
- ▶ Distributed multi-level map
- ▶ CAP: strong consistency and partition tolerance

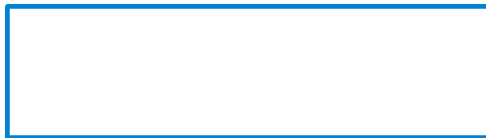


# Data Model



## Data Model (1/5)

- ▶ Table
- ▶ Distributed multi-dimensional sparse `map`





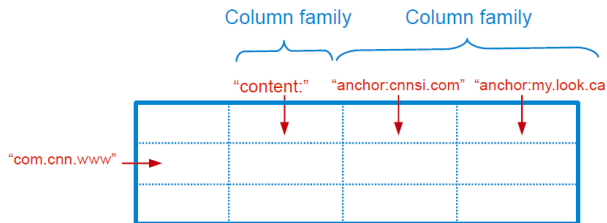
## Data Model (2/5)

- ▶ Rows
- ▶ Every read or write in a row is atomic.
- ▶ Rows sorted in lexicographical order.



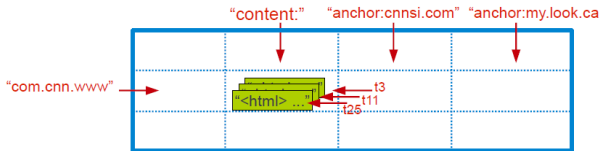
# Data Model (3/5)

- ▶ Column
- ▶ The **basic unit** of data access.
- ▶ **Column families**: group of (the same type) column keys.
- ▶ Column key naming: **family:qualifier**



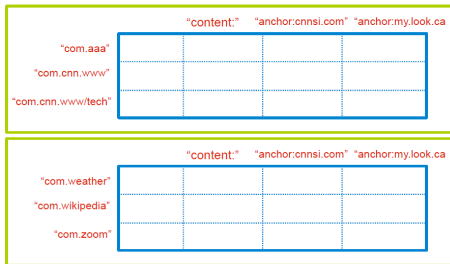
# Data Model (4/5)

- ▶ Timestamp
- ▶ Each column value may contain multiple **versions**.



## Data Model (5/5)

- ▶ **Tablet:** contiguous ranges of rows stored together.
- ▶ Tablets are split by the system when they become too large.
- ▶ Each tablet is served by exactly one tablet server.

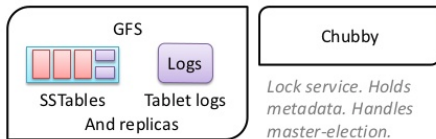
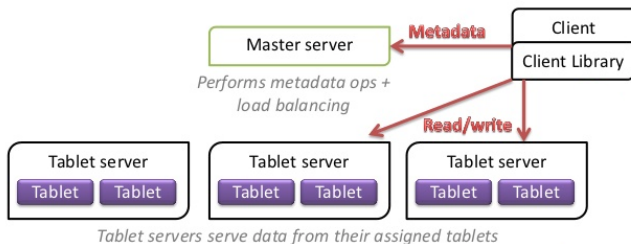






# System Architecture

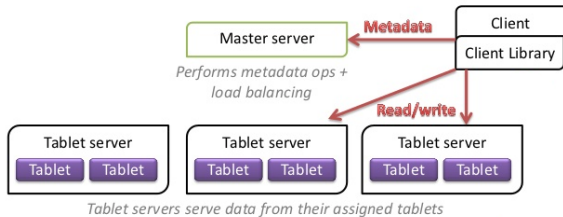
# BigTable System Structure



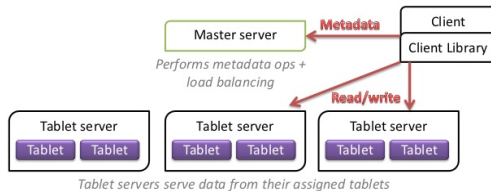
[<https://www.slideshare.net/GrishaWeintraub/cap-28353551>]

# Main Components

- ▶ Master
- ▶ Tablet server
- ▶ Client library

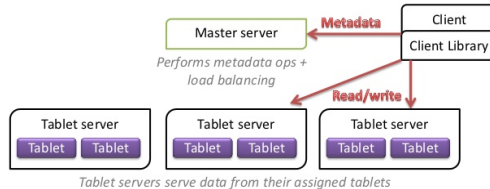


- ▶ Assigns tablets to tablet server.



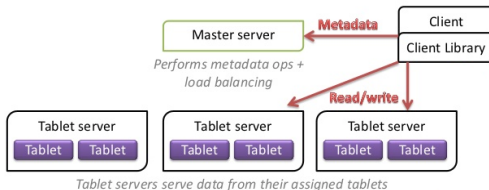
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- ▶ Assigns tablets to tablet server.
- ▶ Balances tablet server load.



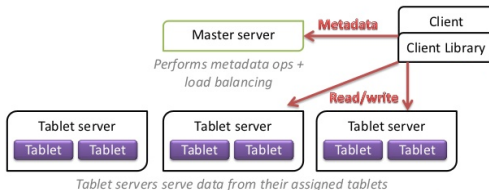
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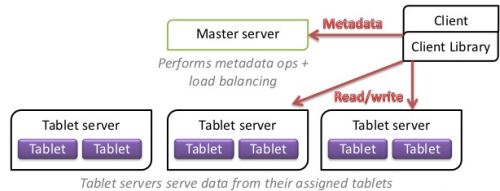
# Master

- ▶ Assigns tablets to tablet server.
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- ▶ Garbage collection of unneeded files in GFS.
- ▶ Handles **schema changes**, e.g., table and column family creations



# Tablet Server

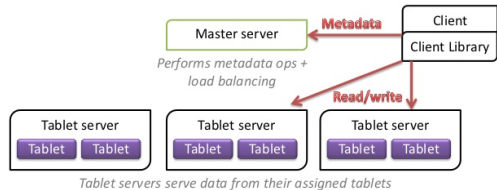
- ▶ Can be added or removed dynamically.





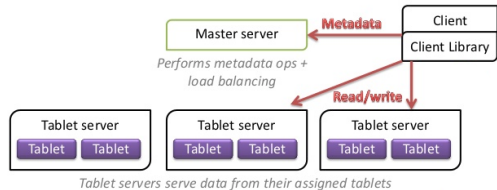
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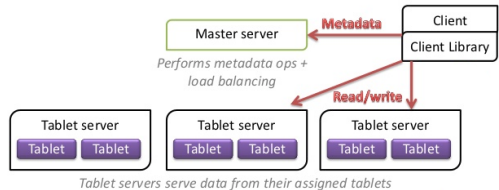
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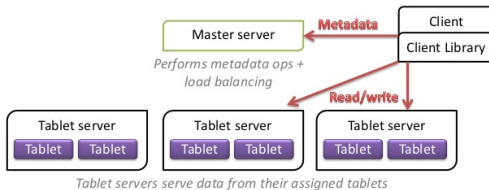
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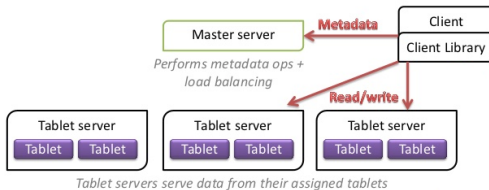
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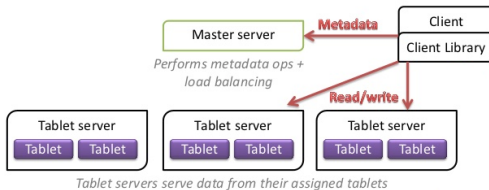
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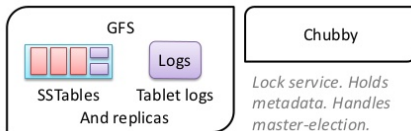
# Client Library

- ▶ Library that is linked into every client.
- ▶ Client **data** does not move through the **master**.
- ▶ Clients communicate **directly** with **tablet servers** for **reads/writes**.



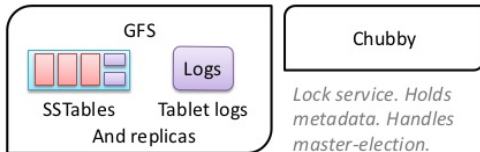
# Building Blocks

- ▶ The building blocks for the BigTable are:
  - Google File System (GFS)
  - Chubby
  - SSTable



# Google File System (GFS)

- ▶ Large-scale distributed file system.
- ▶ Store log and data files.

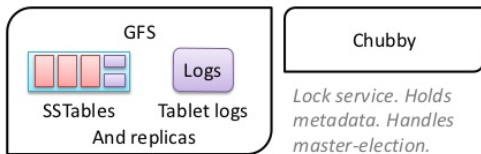






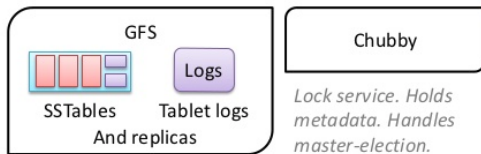
# Chubby Lock Service

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## Chubby Lock Service

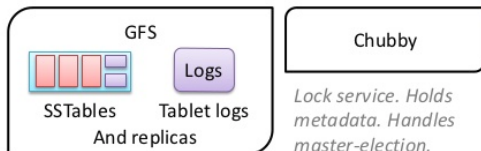
- ▶ Ensure there is only **one active master**.
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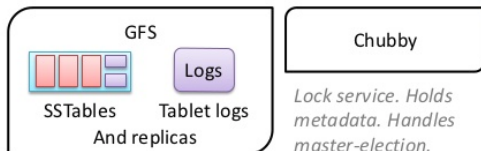
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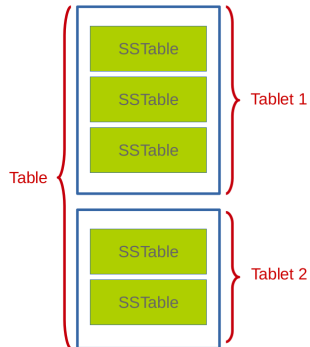
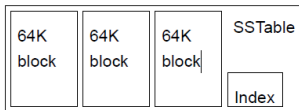
## Chubby Lock Service

- ▶ Ensure there is only **one active master**.
- ▶ Store **bootstrap location** of BigTable data.
- ▶ **Discover** tablet servers.
- ▶ Store BigTable **schema** information and **access control lists**.



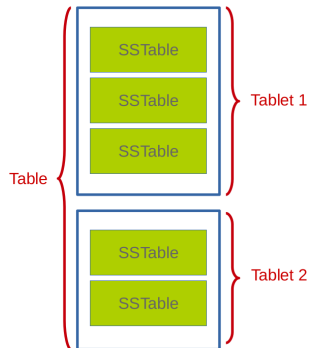
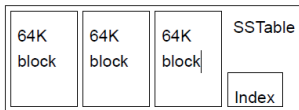
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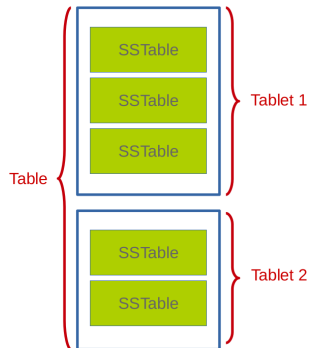
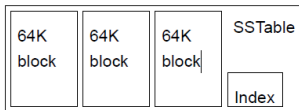
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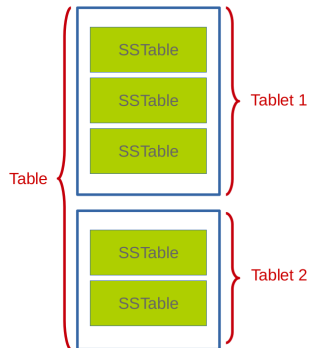
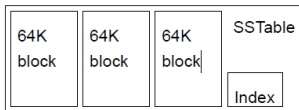
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- ▶ Each SSTable is stored in a **GFS file**.





# Tablet Serving



# Master Startup

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  - **Communicates** with every live tablet server to discover what tablets are already assigned to each server.
  - **Scans the METADATA** table to learn the set of tablets.



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- ▶ Master uses **Chubby** to keep tracks of **live** tablet serves and **unassigned** tablets.
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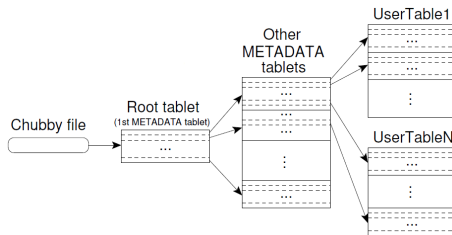


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- ▶ Master detects the **status of the lock of each tablet server** by checking periodically.
- ▶ Master is responsible for finding when tablet server is **no longer serving its tablets** and **reassigning** those tablets as soon as possible.

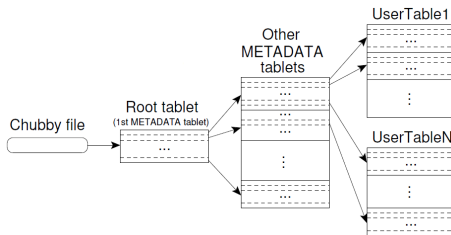
# Finding a Tablet

- ▶ Three-level hierarchy.



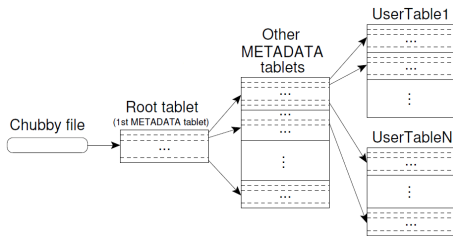
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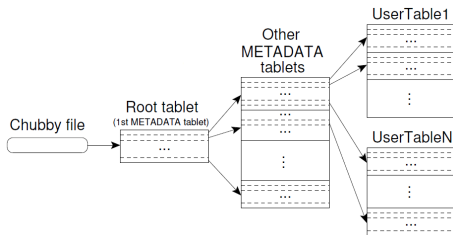
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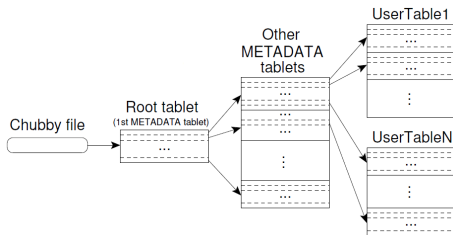
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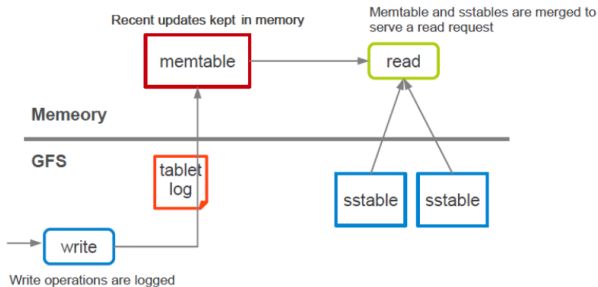
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- ▶ Root tablet contains location of all tablets in a special METADATA table.
- ▶ METADATA table contains location of each tablet under a row.
- ▶ The client library caches tablet locations.



## Tablet Serving (1/2)

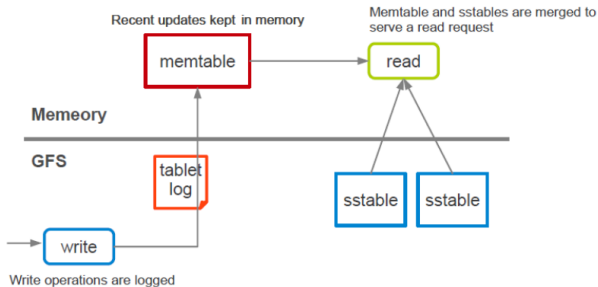
- Updates committed to a **commit log**.





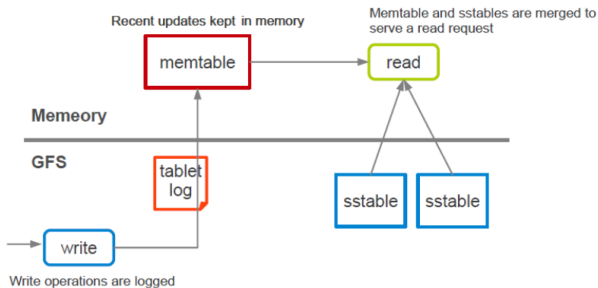
# Tablet Serving (1/2)

- ▶ Updates committed to a **commit log**.
- ▶ Recently committed updates are stored in **memory** - **memtable**



# Tablet Serving (1/2)

- ▶ Updates committed to a **commit log**.
- ▶ Recently committed updates are stored in **memory** - **memtable**
- ▶ **Older updates** are stored in a sequence of **SSTables**.





## Tablet Serving (2/2)

- ▶ Strong consistency
  - Only one tablet server is responsible for a given piece of data.
  - Replication is handled on the GFS layer.



## Tablet Serving (2/2)

- ▶ Strong consistency
  - Only one tablet server is responsible for a given piece of data.
  - Replication is handled on the GFS layer.
- ▶ Trade-off with availability
  - If a tablet server fails, its portion of data is temporarily unavailable until a new server is assigned.



## BigTable vs. HBase

BigTable	HBase
GFS	HDFS
Tablet Server	Region Server
SSTable	StoreFile
Memtable	MemStore
Chubby	ZooKeeper



# HBase Example

```
# Create the table "test", with the column family "cf"  
create 'test', 'cf'
```



# HBase Example

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# Use describe to get the description of the "test" table  
describe 'test'
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# Put data in the "test" table  
put 'test', 'row1', 'cf:a', 'value1'  
put 'test', 'row2', 'cf:b', 'value2'  
put 'test', 'row3', 'cf:c', 'value3'
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# Scan the table for all data at once  
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```

```
# Scan the table for all data at once  
scan 'test'
```

```
# To get a single row of data at a time, use the get command  
get 'test', 'row1'
```



# Cassandra



# Cassandra

- ▶ A **column-oriented** database
- ▶ It was created for **Facebook** and was later **open sourced**
- ▶ **CAP**: **availability** and **partition tolerance**





## Borrowed From BigTable

- ▶ Data model: **column oriented**
  - **Keyspaces** (similar to the schema in a relational database), **tables**, and **columns**.

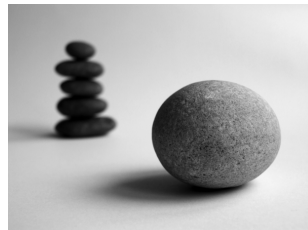


## Borrowed From BigTable

- ▶ Data model: **column oriented**
  - **Keyspaces** (similar to the schema in a relational database), **tables**, and **columns**.
- ▶ **SSTable** disk storage
  - Append-only commit log
  - Memtable (buffering and sorting)
  - Immutable sstable files

## Data Partitioning (1/2)

- ▶ **Key/value**, where values are stored as **objects**.
- ▶ If size of data **exceeds the capacity** of a single machine: **partitioning**



## Data Partitioning (1/2)

- ▶ **Key/value**, where values are stored as **objects**.
- ▶ If size of data **exceeds the capacity** of a single machine: **partitioning**
- ▶ **Consistent hashing** for partitioning.







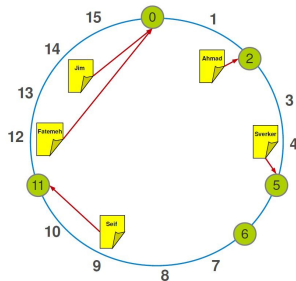
## Data Partitioning (2/2)

- ▶ Consistent hashing.
- ▶ Hash both data and node ids using the same hash function in a same id space.
- ▶  $\text{partition} = \text{hash}(d) \bmod n$ ,  $d$ : data,  $n$ : the size of the id space

## Data Partitioning (2/2)

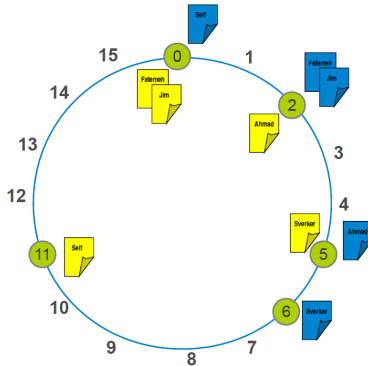
- ▶ Consistent hashing.
- ▶ Hash both data and node ids using the same hash function in a same id space.
- ▶  $\text{partition} = \text{hash}(d) \bmod n$ ,  $d$ : data,  $n$ : the size of the id space

```
id space = [0, 15], n = 16  
hash("Fatemeh") = 12  
hash("Ahmad") = 2  
hash("Seif") = 9  
hash("Jim") = 14  
hash("Sverker") = 4
```



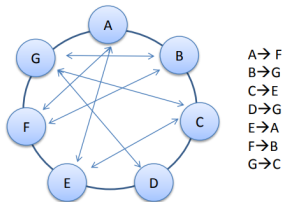
# Replication

- ▶ To achieve high **availability** and **durability**, data should be **replicated** on multiple nodes.



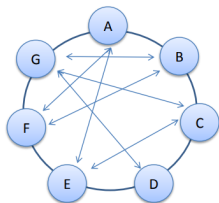
## Adding and Removing Nodes

- ▶ Gossip-based mechanism: **periodically**, each node contacts another randomly selected node.

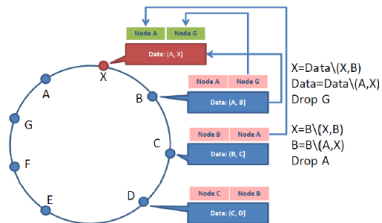


# Adding and Removing Nodes

- Gossip-based mechanism: **periodically**, each node contacts another randomly selected node.



$A \rightarrow F$   
 $B \rightarrow G$   
 $C \rightarrow E$   
 $D \rightarrow G$   
 $E \rightarrow A$   
 $F \rightarrow B$   
 $G \rightarrow C$





# Cassandra Example

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with replication = {'class': 'SimpleStrategy', 'replication_factor': 1};
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# Navigate to the "test" keyspace  
use test
```





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```

```
# Insert a row  
insert into words(word, count) values('hello', 5);
```

```
# Look at the table  
select * from words;
```

# Neo4j



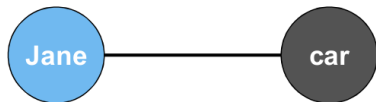
# Neo4j

- ▶ A graph database
- ▶ The relationships between data is equally important as the data itself
- ▶ Cypher: a declarative query language similar to SQL, but optimized for graphs
- ▶ CAP: strong consistency and availability



## ► Node (Vertex)

- The **main data element** from which graphs are constructed.
- A waypoint along a **traversal route**



- ▶ Relationship (Edge)
- ▶ May contain
  - Direction
  - Metadata, e.g., weight or relationship type



## ► Label

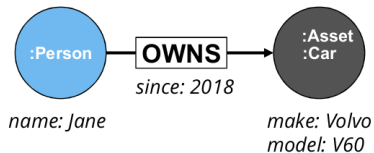
- Define **node category** (optional)
- Can have **more than one**



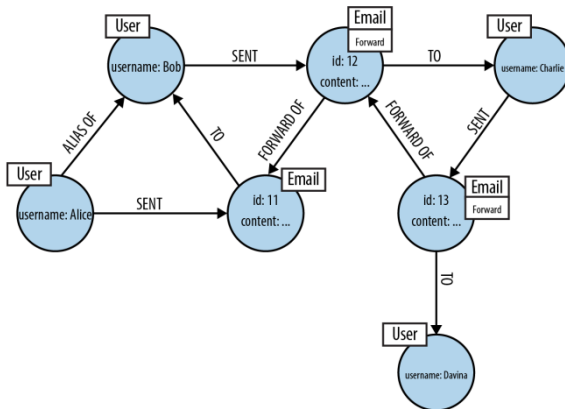


## ► Properties

- **Enrich** a node or relationship



# Example



[Ian Robinson et al., Graph Databases, 2015]

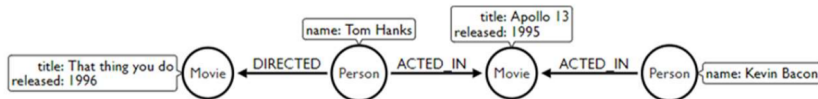




# What is Cypher?

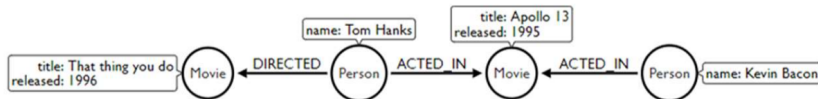
- ▶ Declarative query language
- ▶ `()`: Nodes
- ▶ `[]`: Relationships
- ▶ `{}`: Properties

## Cypher Example (1/2)



```
// Match all nodes  
MATCH (n)  
RETURN n;
```

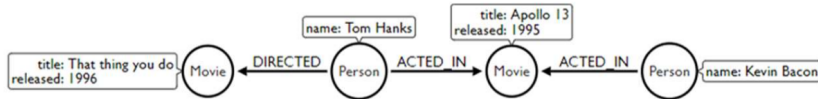
## Cypher Example (1/2)



```
// Match all nodes  
MATCH (n)  
RETURN n;
```

```
// Match all nodes with a Person label  
MATCH (n:Person)  
RETURN n;
```

# Cypher Example (1/2)

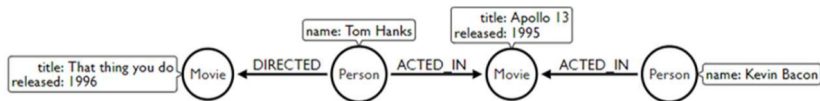


```
// Match all nodes
MATCH (n)
RETURN n;
```

```
// Match all nodes with a Person label
MATCH (n:Person)
RETURN n;
```

```
// Match all nodes with a Person label and property name is 'Tom Hanks'
MATCH (n:Person {name: 'Tom Hanks'})
RETURN n;
```

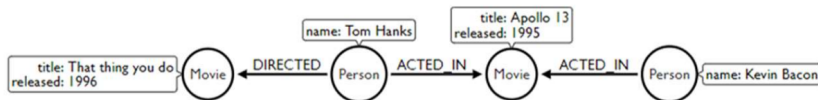
## Cypher Example (2/2)



```
// Return nodes with label Person and name property equals 'Tom Hanks'
MATCH (p:Person)
WHERE p.name = 'Tom Hanks'
RETURN p;
```



## Cypher Example (2/2)



```

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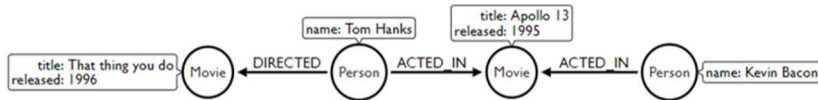
```

```

// Return nodes with label Movie, released property is between 1991 and 1999
MATCH (m:Movie)
WHERE m.released > 1990 AND m.released < 2000
RETURN m;

```

## Cypher Example (2/2)



```

// Return nodes with label Person and name property equals 'Tom Hanks'
MATCH (p:Person)
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RETURN p;

```

```

// Return nodes with label Movie, released property is between 1991 and 1999
MATCH (m:Movie)
WHERE m.released > 1990 AND m.released < 2000
RETURN m;

```

```

// Find all the movies Tom Hanks acted in
MATCH (:Person {name:'Tom Hanks'})-[:ACTED_IN]->(m:Movie)
RETURN m.title;

```

# Summary



## Summary

- ▶ NoSQL data models: key-value, column-oriented, document-oriented, graph-based
- ▶ CAP (Consistency vs. Availability)



## Summary

- ▶ BigTable
- ▶ Column-oriented
- ▶ Main components: master, tablet server, client library
- ▶ Basic components: GFS, SSTable, Chubby
- ▶ CP



## Summary

- ▶ Cassandra
- ▶ Column-oriented (similar to BigTable)
- ▶ Consistency hashing
- ▶ Gossip-based membership
- ▶ AP



## Summary

- ▶ Neo4j
- ▶ Graph-based
- ▶ Cypher
- ▶ CA



## References

- ▶ F. Chang et al., Bigtable: A distributed storage system for structured data, ACM Transactions on Computer Systems (TOCS) 26.2, 2008.
- ▶ A. Lakshman et al., Cassandra: a decentralized structured storage system, ACM SIGOPS Operating Systems Review 44.2, 2010.
- ▶ I. Robinson et al., Graph Databases (2nd ed.), O'Reilly Media, 2015.



# Questions?

## Acknowledgements

Some content of the Neo4j slides were derived from Ljubica Lazarevic's slides.