# Large Scale Graph Processing 

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https://id2221kth.github.io

## https://tinyurl.com/hk7hzpw5

## Where Are We?



- A flexible abstraction for describing relationships between discrete objects.




## Graph Algorithms Challenges

- Difficult to extract parallelism based on partitioning of the data.
- Difficult to express parallelism based on partitioning of computation.


## Graph Algorithms Challenges

- Difficult to extract parallelism based on partitioning of the data.
- Difficult to express parallelism based on partitioning of computation.
- Graph partition is a challenging problem.
- Partition large scale graphs and distribut to hosts.



## Edge-Cut Graph Partitioning

- Divide vertices of a graph into disjoint clusters.
- Nearly equal size (w.r.t. the number of vertices).
- With the minimum number of edges that span separated clusters.



## Vertex-Cut Graph Partitioning

- Divide edges of a graph into disjoint clusters.
- Nearly equal size (w.r.t. the number of edges).
- With the minimum number of replicated vertices.



## Edge-Cut vs. Vertex-Cut Graph Partitioning (1/2)

- Natural graphs: skewed Power-Law degree distribution.
- Edge-cut algorithms perform poorly on Power-Law Graphs.



Edge-Cut vs. Vertex-Cut Graph Partitioning (2/2)


## PageRank with MapReduce



$$
R[i]=\sum_{j \in \operatorname{Nbrs}(i)} w_{j i} R[j]
$$

(Kxit PageRank Example (1/2)

- $R[i]=\sum_{j \in \operatorname{Mbrs}(i)} W_{j i} R[j]$



## K표 PageRank Example (1/2)

- $R[i]=\sum_{j \in \operatorname{Mbrs}(i)} w_{j i} R[j]$
- Input


```
V1: [0.25, V2, V3, V4]
V2: [0.25, V3, V4]
V3: [0.25, V1]
V4: [0.25, V1, V3]
```


## 톤

- $R[i]=\sum_{j \in \operatorname{Mbrs}(i)} W_{j i} R[j]$
- Input


```
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V2: [0.25, V3, V4]
V3: [0.25, V1]
V4: [0.25, V1, V3]
```

- Share the rank among all outgoing links

```
V1: (V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3)
V2: (V3, 0.25/2), (V4, 0.25/2)
V3: (V1, 0.25/1)
V4: (V1, 0.25/2), (V3, 0.25/2)
```


## 톤 . PageRank Example (2/2)

- $R[i]=\sum_{j \in \operatorname{Mbrs}(i)} W_{j \mathrm{i}} R[j]$


```
V1: (V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3)
V2: (V3, 0.25/2), (V4, 0.25/2)
V3: (V1, 0.25/1)
V4: (V1, 0.25/2), (V3, 0.25/2)
```


## PageRank Example (2/2)

- R[i] $=\sum_{j \in \operatorname{Nbrs}(i)} W_{j i} R[j]$


```
V1: (V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3)
V2: (V3, 0.25/2), (V4, 0.25/2)
V3: (V1, 0.25/1)
V4: (V1, 0.25/2), (V3, 0.25/2)
```

- Output after one iteration

```
V1: [0.37, V2, V3, V4]
V2: [0.08, V3, V4]
V3: [0.33, V1]
V4: [0.20, V1, V3]
```


## PageRank in MapReduce - Map (1/2)

- Map function


```
map(key: [url, pagerank], value: outlink_list)
    for each outlink in outlink_list:
            emit(key: outlink, value: pagerank / size(outlink_list))
    emit(key: url, value: outlink_list)
```


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```

- Input (key, value)

```
((V1, 0.25), [V2, V3, V4])
((V2, 0.25), [V3, V4])
((V3, 0.25), [V1])
((V4, 0.25), [V1, V3])
```


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    emit(key: url, value: outlink_list)
```

- Intermediate (key, value)

```
(V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3), (V3, 0.25/2), (V4, 0.25/2), (V1, 0.25/1),
(V1, 0.25/2), (V3, 0.25/2)
(V1, [V2, V3, V4])
(V2, [V3, V4])
(V3, [V1])
(V4, [V1, V3])
```


## PageRank in MapReduce - Shuffle

- Intermediate (key, value)

```
(V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3), (V3, 0.25/2), (V4, 0.25/2), (V1, 0.25/1),
(V1, 0.25/2), (V3, 0.25/2)
(V1, [V2, V3, V4])
(V2, [V3, V4])
(V3, [V1])
(V4, [V1, V3])
```


## PageRank in MapReduce - Shuffle

- Intermediate (key, value)

```
(V2, 0.25/3), (V3, 0.25/3), (V4, 0.25/3), (V3, 0.25/2), (V4, 0.25/2), (V1, 0.25/1),
(V1, 0.25/2), (V3, 0.25/2)
(V1, [V2, V3, V4])
(V2, [V3, V4])
(V3, [V1])
(V4, [V1, V3])
```

- After shuffling

```
(V1, 0.25/1), (V1, 0.25/2), (V1, [V2, V3, V4])
(V2, 0.25/3), (V2, [V3, V4])
(V3, 0.25/3), (V3, 0.25/2), (V3, 0.25/2), (V3, [V1])
(V4, 0.25/3), (V4, 0.25/2), (V4, [V1, V3])
```


## PageRank in MapReduce - Reduce (1/2)

- Reduce function

```
reducer(key: url, value: list_pr_or_urls)
    outlink_list = []
    pagerank = 0
    for each pr_or_urls in list_pr_or_urls:
        if is_list(pr_or_urls):
            outlink_list = pr_or_urls
        else
            pagerank += pr_or_urls
    emit(key: [url, pagerank], value: outlink_list)
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```

- Input of the Reduce function

```
(V1, 0.25/1), (V1, 0.25/2), (V1, [V2, V3, V4])
(V2, 0.25/3), (V2, [V3, V4])
(V3, 0.25/3), (V3, 0.25/2), (V3, 0.25/2), (V3, [V1])
(V4, 0.25/3), (V4, 0.25/2), (V4, [V1, V3])
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- Output

```
((V1, 0.37), [V2, V3, V4])
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## Problems with MapReduce for Graph Analytics

- MapReduce does not directly support iterative algorithms.
- Invariant graph-topology-data re-loaded and re-processed at each iteration is wasting I/O, network bandwidth, and CPU


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- MapReduce does not directly support iterative algorithms.
- Invariant graph-topology-data re-loaded and re-processed at each iteration is wasting I/O, network bandwidth, and CPU
- Materializations of intermediate results at every MapReduce iteration harm performance.


## Think Like a Vertex

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- Each vertex computes individually its value (in parallel).
- Computation typically depends on the neighbors.



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- Each vertex computes individually its value (in parallel).
- Computation typically depends on the neighbors.
- Also know as graph-parallel processing model.


Data-Parallel vs. Graph-Parallel Computation


Graph-Parallel

"ctuas

Pregel

## Pregel

- Large-scale graph-parallel processing platform developed at Google.
- Inspired by bulk synchronous parallel (BSP) model.



## Execution Model (1/2)

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- A vertex in superstep S can:
- reads messages sent to it in superstep $S-1$.
- sends messages to other vertices: receiving at superstep $\mathrm{S}+1$.
- modifies its state.


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- Applications run in sequence of iterations, called supersteps.
- A vertex in superstep S can:
- reads messages sent to it in superstep $S-1$.
- sends messages to other vertices: receiving at superstep $\mathrm{S}+1$.
- modifies its state.
- Vertices communicate directly with one another by sending messages.


## Execution Model (2/2)

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- Superstep 0: all vertices are in the active state.
- A vertex deactivates itself by voting to halt: no further work to do.
- A halted vertex can be active if it receives a message.
- The whole algorithm terminates when:
- All vertices are simultaneously inactive.
- There are no messages in transit.



## Example: Max Value (1/4)

```
i_val := val
for each message m
    if m > val then val := m
if i_val == val then
    vote_to_halt
else
    for each neighbor v
        send_message(v, val)
```



## Example: Max Value (2/4)

```
i_val := val
for each message m
    if m > val then val := m
if i_val == val then
    vote_to_halt
else
    for each neighbor v
        send_message(v, val)
```



## Example: Max Value (3/4)

```
i_val := val
for each message m
    if m > val then val := m
if i_val == val then
    vote_to_halt
else
    for each neighbor v
        send_message(v, val)
```



Example: Max Value (4/4)

```
i_val := val
for each message m
    if m > val then val := m
if i_val == val then
    vote_to_halt
else
    for each neighbor v
        send_message(v, val)
```



## Example: PageRank



$$
R[i]=\sum_{j \in \operatorname{Nbrs}(i)} W_{j i} R[j]
$$

## Example: PageRank

```
Pregel_PageRank(i, messages):
    // receive all the messages
    total = 0
    foreach(msg in messages):
        total = total + msg
    // update the rank of this vertex
    R[i] = total
    // send new messages to neighbors
    foreach(j in out_neighbors[i]):
        sendmsg(R[i] * wij) to vertex j
```

$$
R[i]=\sum_{j \in \operatorname{Nbrs}(i)} W_{j i} R[j]
$$

## GraphLab/Turi

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- GraphLab allows asynchronous iterative computation.
- Vertex scope of vertex v: the data stored in v, and in all adjacent vertices and edges.
- A vertex can read and modify any of the data in its scope (shared memory).



## Example: PageRank (GraphLab)

```
GraphLab_PageRank(i)
    // compute sum over neighbors
    total = 0
    foreach(j in in_neighbors(i)):
        total = total + R[j] * wji
    // update the PageRank
    R[i] = total
    // trigger neighbors to run again
    foreach(j in out_neighbors(i)):
        signal vertex-program on j
```

$$
R[i]=\sum_{j \in \operatorname{Nbrs}(i)} W_{j i} R[j]
$$

## Gather-Apply-Scatter (GAS)

- Factorizes the local vertices functions into the Gather, Apply and Scatter phases.
- Gather: accumulate information from neighborhood.
- Apply: apply the accumulated value to center vertex.
- Scatter: update adjacent edges and vertices.


## Example: PageRank (GraphLab - GAS)

```
PowerGraph_PageRank(i):
    Gather(j -> i):
        return wji * R[j]
    sum(a, b):
        return a + b
    // total: Gather and sum
    Apply(i, total):
        R[i] = total
```

    Scatter (i \(->\) j):
        if \(R[i]\) changed then activate( \(j\) )
    $$
R[i]=\sum_{j \in \operatorname{Nbrs}(i)} W_{j i} R[j]
$$

## Think Like a Table

Data-Parallel vs. Graph-Parallel Computation


Graph-Parallel


- Graph-parallel computation: restricting the types of computation to achieve performance.


## Motivation (2/3)

- Graph-parallel computation: restricting the types of computation to achieve performance.
- The same restrictions make it difficult and inefficient to express many stages in a typical graph-analytics pipeline.


Live-Journal: 69 Million Edges


## Motivation (3/3)




## Think Like a Table

- Unifies data-parallel and graph-parallel systems.
- Tables and Graphs are composable views of the same physical data.



GraphX Unified Representation


Graph View

## GraphX

## GraphX

- GraphX is the library to perform graph-parallel processing in Spark.



## The Property Graph Data Model

- Spark represent graph structured data as a property graph.
- It is logically represented as a pair of vertex and edge property collections.
- VertexRDD and EdgeRDD

```
// VD: the type of the vertex attribute
// ED: the type of the edge attribute
class Graph[VD, ED] {
    val vertices: VertexRDD[VD]
    val edges: EdgeRDD[ED]
}
```



| Id | Property ( V ) |  |
| :---: | :---: | :---: |
| 3 | (rxin, student) |  |
| 7 | (jgonzal, postdoc) |  |
| 5 | (franklin, professor) |  |
| 2 | (istoica, professor) |  |
| Edge Table |  |  |
| Srcld | Dstid | Property (E) |
| 3 | 7 | Collaborator |
| 5 | 3 | Advisor |
| 2 | 5 | Colleague |
| 5 | 7 | PI |

## The Vertex Collection

- VertexRDD: contains the vertex properties keyed by the vertex ID.

```
class Graph[VD, ED] {
    val vertices: VertexRDD [VD]
    val edges: EdgeRDD [ED]
}
```

// VD: the type of the vertex attribute abstract class VertexRDD[VD] extends RDD[(VertexId, VD)]


## The Edge Collection

- EdgeRDD: contains the edge properties keyed by the source and destination vertex IDs.

```
class Graph[VD, ED] {
    val vertices: VertexRDD[VD]
    val edges: EdgeRDD[ED]
}
// ED: the type of the edge attribute
case class Edge[ED] (srcId: VertexId, dstId: VertexId, attr: ED)
abstract class EdgeRDD[ED] extends RDD[Edge[ED]]
```



VertexTable


Edge Table

| Sreld | Dstld | Property (E) |
| :---: | :---: | :---: |
| 7 | 7 | Colibocotar |
| 5 | 3 | Alvor |
| 2 | 5 | Collegue |
| 5 | 7 | P1 |
| 5 | 7 |  |

Edges: (A-a-B

## The Triplet Collection

- The triplets collection consists of each edge and its corresponding source and destination vertex properties.



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- It logically joins the vertex and edge properties: RDD [EdgeTriplet [VD, ED] ].



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- The triplets collection consists of each edge and its corresponding source and destination vertex properties.
- It logically joins the vertex and edge properties: RDD [EdgeTriplet [VD, ED]].
- The EdgeTriplet class extends the Edge class by adding the srcAttr and dstAttr members, which contain the source and destination properties respectively.



## Building a Property Graph


val users: RDD[(VertexId, (String, String))] = sc.parallelize(Array((3L, ("rxin", "student")), (7L, ("jgonzal", "postdoc")), (5L, ("franklin", "prof")), (2L, ("istoica", "prof"))))

## Building a Property Graph


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val relationships: RDD[Edge[String]] = sc.parallelize(Array(Edge(3L, 7L, "collab"), Edge(5L, 3L, "advisor"), Edge(2L, 5L, "colleague"), Edge(5L, 7L, "pi"), Edge(5L, 1L, "-")))

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val defaultUser = ("John Doe", "Missing")

## Building a Property Graph



| Vertex Table |  |  |
| :---: | :---: | :---: |
| $1 d$ | Property (V) |  |
| 3 | (rxin, student) |  |
| 7 | (gonzal postdoc) |  |
| 5 | (franklin, professor) |  |
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| Edge Table |  |  |
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val users: RDD[(VertexId, (String, String))] = sc.parallelize(Array((3L, ("rxin", "student")),
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    Edge(5L, 3L, "advisor"), Edge(2L, 5L, "colleague"), Edge(5L, 7L, "pi"), Edge(5L, 1L, "-")))
```

val defaultUser = ("John Doe", "Missing")
val graph: Graph[(String, String), String] = Graph(users, relationships, defaultUser)

- Information about the graph
- Property operators
- Structural operators
- Joins
- Aggregation
- Iterative computation
- ...


## Summary

## Summary

- Think like a vertex
- Pregel: BSP, synchronous parallel model, message passing
- GraphLab: asynchronous model, shared memory, GAS
- Think like a table
- Graphx: unifies data-parallel and graph-parallel systems.


## References

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- J. Gonzalez et al., "Powergraph: distributed graph-parallel computation on natural graphs", OSDI 2012
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## Questions?

